



**UNIVERSITI PUTRA MALAYSIA**

**PROACTIVE TRAFFIC-ADAPTIVE TUNING OF CONTENTION WINDOW  
FOR WIRELESS SENSOR NETWORK MEDIUM-ACCESS CONTROL  
PROTOCOL**

**NESAE MOUZEHKESH PIRBORJ**

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PROTOCOL**

**By**

**NESAE MOUZEHKESH PIRBORJ**

**Thesis Submitted to the School of Graduate Studies, Universiti Putra Malaysia, in Fulfilment of  
the Requirements for the Degree of Master of Science**

**December 2009**



## **DEDICATION**

This thesis is dedicated to my parents and all those whom I have ever loved and will

love...

God bless you all...



Abstract of thesis presented to the Senate of Universiti Putra Malaysia in fulfilment of the requirement for the degree of Master of Science

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**December 2009**

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The ongoing advances in wireless networks have further expanded the boundaries to the new and challenging area of Wireless Sensor Networks (WSN). Unique properties of sensor nodes such as limited energy storage, constrained processing capabilities and the especially different environments they are usually deployed in have prompted the need of novel protocols in all the layers of the communication stack. A Medium Access Control (MAC) protocol is responsible to sufficiently provide access to a shared medium. Therefore effective techniques in order to reduce the probability of collisions while contending for the medium can be established in a MAC protocol for it organizes the specific time slot a node can have access to the channel. The need for further improving the current applied MAC protocols for WSN in order to reduce the probability of collisions while being energy aware has motivated this research. Sensor MAC as the very first MAC protocol for WSN has been designed on top of the IEEE 802.11 MAC protocol along with some added features to meet the special requirements of a WSN. However the Back-Off scheme of Sensor MAC (S-MAC) is based on a fixed Contention Window (CW) size. This is



known as a significant trouble spot in S-MAC in the sense that the delay produced during collisions and idle listening can be so critical to the limited battery lifetime of a sensor node. IEEE 802.11 MAC protocol follows a static approach for obtaining the back-off time and resets the CW to its default minimum upon just one successful transmission and doubles it each time it faces a collision. While the back-off algorithm of IEEE 802.11 suffers from unfairness for its faulty behaviour in both high and low traffic loads the back-off mechanism in S-MAC suffers from a fixed CW size. Reducing the undesired idle listening time caused by unnecessary long back-off times when traffic is low and also decreasing the probability of collisions in situations with high traffic load due to the fixed CW size in S-MAC have motivated our research. We have tried to come up with a dynamic back-off algorithm for S-MAC that can extract the current traffic information of the network and engage them in estimating the contention window from which the back-off time is chosen. Our approach is a proactive algorithm to get the CW of the neighbouring nodes ready before contending for the medium. The performance of our algorithm has been measured in terms of average delay, average throughput, delivery ratio, and average energy efficiency. It is shown that our back-off scheme has reduced the delay by 47% and has decreased the energy consumption up to above 15% over the current S-MAC implementation. The delivery ratio and throughput have been improved up to 44% and 28% respectively.

Abstrak tesis yang dikemukakan kepada Senat Universiti Putra Malaysia  
sebagai memenuhi keperluan untuk ijazah Master of Sains

**PENUMPUAN PENALA TRAFIK-TETINGKAP PROAKTIF PENYATA  
UNTUK RANGKAIAN PENDERIAAN TANDA WAYAR PROTOKOL  
KAWALAN AKSES SEDERHANA**

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Kemajuan yang berterusan di dalam bidang rangkaian tanpa wayar telah mengembangkan sempadannya menuju kepada sebuah bidang yang mencabar rangkaian Penderia Tanpa Wayar (WSN). Nod penderia yang unik seperti penyimpanan tenaga terhad, kekangan keupayaan pemprosesan terutamanya di peskitaran yang berbeza dimana nod-nod ini kebiasaannya diletakkan telah mendorong keperluan protokol-protokol novel di dalam semua lapisan komunikasi. Medium Access Control (MAC) bertanggungjawab untuk menyediakan akses yang cukup kepada perkongsian medium. Ini bermakna, teknik yang berkesan untuk mengurangkan kebarangkalian pelanggaran semasa bertanding untuk pengantara boleh dibangunkan dalam sebuah protokol mac untuk ia mengatur slot masa tertentu sesebuah nod yang mempunyai akses kepada saluran. Keperluan untuk menambah baik penggunaan protocol MAC mengurangkan kebarangkalian pelanggaran dan pada waktu yang sama telah meransang kajian ini. Penderia MAC sebagai protocol MAC yang pertama untuk WSN telah di reka bentuk di atas protocol IEEE 802.11 bersama dengan penambahan ciri untuk memenuhi keperluan sesebuah WSN. Walaubagaimana pun, skim Back-OFF S-MAC adalah berdasarkan saiz Contention



Window yang tetap. Ini dikenali sebagai masalah dalam S-MAC yang paling ketara dalam konteks kelewatannya yang dihasilkan semasa pelanggaran dan pendengaran serta boleh menjadi sangat kritikal kepada jangka hayat bateri yang terhad dalam sesebuah nod penerima. Protokol MAC IEEE 802.11 mengikuti pendekatan statik untuk mendapatkan masa back-off dan menetapkan semula CW kepada nilai lalai minimum setelah satu penghantaran berjaya dan menggandakannya setiap kali berlakunya pelanggaran. Sementara algorithm back-off mengalami kegagalan di dalam kedua-dua trafik (trafik rendah dan tinggi), kekurangan mekanisma back-off dalam S-MAC adalah daripada saiz cw yang tetap. Pengurangan masa pendengaran disebabkan oleh masa Back-Off yang panjang dan tidak diperlukan semasa trafik rendah dan juga mengurangkan kebarangkalian pelanggaran dalam situasi-situasi trafik beban tinggi oleh sebab saiz cw yang tetap dalam s-mac telah merangsang kajian kami. Kami telah mencuba untuk menghasilkan algorithm back-off dinamik untuk s-mac yang boleh mengestrak maklumat trafik semasa jaringan dan menggabungkan mereka dalam menganggar pernyataan tettingkap dari mana masa back-off dipilih. Pendekatan kami adalah menghasilkan sebuah algoritma proaktif untuk mendapatkan cw nod-nod yang bersebelahan bersedia sebelum bertanding untuk pengantara. Prestasi algoritma ini telah di ukur dalam konteks kelewatan purata, throughput purata, nisbah penghantaran dan purata keberkesanan tenaga. Ini menunjukkan skim back-off ini telah mengurangkan kelewatan sehingga 47% dan telah mengurangkan penggunaan tenaga sehingga melepasi 15% daripada pelaksanaan S-MAC semasa. Nisbah penghantaran dan throughput telah ditingkatkan sehingga 44% dan 28% masing-masing.

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## **APPROVAL**

I certify that an Examination Committee has met on ..... to conduct the final examination of Nesae Mouzehkesh Pirborj on her Master of Science thesis “Proactive Traffic Adaptive Tuning of Contention Window in Medium Access Control Protocol with Energy Conservation for Wireless Sensor Network” in accordance with Universiti Putra Malaysia (Higher Degree) Act 1980 and Universiti Pertanian Malaysia (Higher Degree) Regulations 1981. The Committee recommends that the candidate be awarded the relevant degree.

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## **DECLARATION**

I hereby declare that the thesis is based on my original work except for quotations and citations which have been duly acknowledged. I also declare that it has not been previously or concurrently submitted for any other degree at UPM or other institutions.

---

**NESAE MOUZEHKESH PIRBORJ**

Date:

## TABLE OF CONTENTS

	<b>Page</b>
<b>DEDICATION</b>	ii
<b>ABSTRACT</b>	iii
<b>ABSTRAK</b>	iv
<b>ACKNOWLEDGEMENT</b>	vii
<b>APROVAL</b>	x
<b>DECLARATION</b>	xi
<b>LIST OF TABLES</b>	xiv
<b>LIST OF FIGURES</b>	xv
<b>LIST OF ABBREVIATIONS</b>	xvii
<b>LIST OF SYMBOLS</b>	xx
<b>1 INTRODUCTION</b>	
1.1 Background	1
1.2 Problem Statement and Motivation	2
1.3 Research Aim and Objectives	6
1.4 Thesis Scope	7
1.5 Study Module	8
1.6 Research Contribution	10
1.7 Thesis Organization	10
<b>2 LITERATURE REVIEW</b>	
2.1 Background	12
2.2 Wireless Sensor Networks	13
2.3 Wireless Sensor Network Protocol Stack	16
2.4 Medium Access Control Protocols	17
2.4.1 Contention-Free MAC Protocols	19
2.4.2 Contention-Based MAC Protocols	21
2.5 MAC Design Space for WSN	23
2.5.1 Energy Efficiency	22
2.5.2 Scalability and Adaptability	25
2.5.3 Fairness	26
2.5.4 Latency	27
2.5.5 Throughput	27
2.6 MAC Protocols Classification for WSN	28
2.6.1 Synchronized Low Duty Cycle Protocols	29
2.6.1.1 S-MAC	29
2.6.2 Unsynchronized Low Duty Cycle Protocols	34
2.7 Fixed Contention Window Size in S-MAC	35
2.8 Traffic Adaptive Back-Off algorithms	35
2.9 Summary	41
<b>3 METHODOLOGY</b>	
3.1 Introduction	44
3.2 Methodology in Brief	47
3.2.1 Assumptions	48
3.2.2 Protocol Operation	48



3.3	Proactive Algorithm	52
3.4	Traffic Adaptive Contention Window Adjustment Algorithm	53
3.5	Implementation in NS-2	60
3.6	Performance metrics	62
3.7	Summary	64
<b>4</b>	<b>RESULTS AND DISCUSSION</b>	
4.1	Background	65
4.2	Network Configuration and scenarios	65
4.3	Performance Evaluation	68
	4.3.1 Delay Analysis	71
	4.3.2 Throughput Analysis	75
	4.3.3 Delivery Ratio Analysis	77
	4.3.4 Energy Consumed Analysis	79
4.4	Comparison of Higher Duty Cycles in S-MAC	83
4.5	Summary	88
<b>5</b>	<b>CONCLUSION</b>	
5.1	Conclusion	90
5.2	Thesis Contribution	90
5.3	Future Work	91
	REFERENCES	93
	APPENDIX I (C++ Codes Added to the S-MAC)	98
	APPENDIX II (TCL Code of the Simulation, Linear Topology)	102
	APPENDIX III (TCL Code of the Simulation)	106
	APPENDIX IV (AWK Script)	110
	APPENDIX V (Energy Model in NS-2)	121
	BIODATA OF STUDENT	123
	LIST OF PUBLICATIONS	124



## LIST OF TABLES

<b>Table</b>	<b>Page</b>
Table 2.1: Summary of the most related reviewed literature	43
Table 4.1: Simulation parameters	67
Table 4.2: MAC and PHY layer parameters	67
Table 4.3: Numerical example for delivery ratio	77
Table 4.4: Consumed energy at sink (10 nodes)	82
Table 4.5: Performance Evaluation of Different Duty Cycles in S-MAC and the Modified S-MAC with Duty Cycle of 10%	87
Table 4.6: Measured enhancement	89



## LIST OF FIGURES

<b>Figure</b>	<b>Page</b>
Figure 1.1: BEB algorithm	4
Figure 1.2: Study module	9
Figure 2.1: The structure of a WSN, a sensor field with a number of sensors collaborating to do a certain task	15
Figure 2.2: Simple structure of a sensor node	16
Figure 2.3: Wireless sensor network protocol stack	17
Figure 2.4a: Basic CSMA/CA scheme (2-way handshaking)	19
Figure 2.4b: CSMA/CA with RTS/CTS (4-way handshaking)	19
Figure 2.5: Hidden terminal problem	22
Figure 2.6: Periodic listen-and-sleep schedule in S-MAC	32
Figure 2.7: Listen period of S-MAC	32
Figure 3.1: Process flow diagram	45
Figure 3.2a: Proposed protocol behaviour	49
Figure 3.2b: Nodes with ready CW upon receiving RTS	50
Figure 3.2c: Nodes with ready CW upon receiving CTS	50
Figure 3.3: Proposed protocol operation	52
Figure 3.4: Virtual transmission time	53
Figure 3.5: Elements needed to calculate the CW	58
Figure 3.6: CW adjustment in a virtual transmission time	60
Figure 3.7: File structure in NS-2	62
Figure 4.1a: Network topology (10 nodes)	66
Figure 4.1b: Network topology (20 nodes)	66
Figure 4.2: Ten-hop linear network with one source and one sink	73
Figure 4.3: Average Delay vs. Number of hops (Linear network)	74



Figure 4.4: Average end-to-end delay vs. Message inter arrival rate	74
Figure 4.5: Throughput vs. Message inter arrival rate	76
Figure 4.6: Delivery ratio vs. Message inter arrival rate	78
Figure 4.7: Average energy consumed vs. Message inter arrival rate	81
Figure 4.8: Energy consumed at sink (10 nodes)	78
Figure 4.9a: High duty cycle	83
Figure 4.9b: Low duty cycle	83
Figure 4.10: Average end-to-end delay vs. Message inter arrival rate; effect of high duty cycles	85
Figure 4.11: Delivery ratio vs. Message inter arrival rate; effect of high duty cycles	86
Figure 4.12: Average energy consumed vs. Message inter arrival rate; effect of high duty cycles	87

## LIST OF ABBREVIATIONS

ACK	Acknowledgment
ADC	Analog to Digital Converter
AODV	Ad-hoc On-demand Distance Vector (routing protocol)
AWK	a perl or C++ script to interpret the results ( <b>AWK</b> stands for <b>A</b> ho, <b>W</b> einberger & <b>K</b> ernighan, the authors of this language)
BEB	Binary Exponential Back-off
B-MAC	Berkeley Medium Access Control
BT	Back-off Time
CBR	Constant Bit Rate
CDMA	Code Division Multiple Access
CPU	Computer Processing Unit
CSB	Channel Status Bit
CSMA	Carrier Sense Multiple Access
CSMA/CA	Carrier Sense Multiple Access with Collision Avoidance
CTS	Clear To Send
CW	Contention Window
DCF	Distributed Coordination Function
DIFS	Distributed Coordination Function Inter Frame Space
D-MAC	Data gathering Medium Access Control Protocol
EIED	Exponential Increase Exponential Decrease
FDMA	Frequency Division Multiple Access
GHz	Giga Hertz
Kbps	Kilo byte per second
LAN	Local Area Network



LEACH	Low Energy Adaptive Clustering Hierarchy
MAC	Medium Access Control
MACA	Multiple Access Collision Avoidance
MACAW	Multiple Access Collision Avoidance for Wireless
Mbps	Mega byte per second
MANET	Mobile Ad hoc Network
NAV	Network Allocation Vector
NAVC	Network Allocation Vector Counter
NIC	Network Interface Card
Otcl	Object tool command language
PCF	Point Coordination Function
PHY	Physical
QoS	Quality of Service
RAM	Random Access Memory
RTS	Request To Send
SDMA	Space Division Multiple Access
SIFS	Short Inter Frame Space
S-MAC	Sensor Medium Access control
Std	Standard
SYNC	Synchronization
TCL	Tool Command Language
TDMA	Time Division Multiple Access
T-MAC	Time-out Medium Access Control
WLAN	Wireless Local Area Network
WPAN	Wireless Personal Area Network
WSN	Wireless Sensor Network



Z-MAC      Zebra Medium Access Control



## LIST OF SYMBOLS

$d_{cs}$	Carrier sense delay
$d_{end-to-end}$	End-to-end delay
$d_{proc}$	Processing delay
$d_{prop}$	Propagation delay
$d_{trans}$	Transmission delay
$Coll_i$	$i^{th}$ collision in a virtual transmission time
dBm	Power ratio in decibels
$E[B]$	Average back-off time
$E[B_1]$	Average back-off time for only one active station
$E[Coll]$	Average length of a collision in a virtual transmission time
$E[N_c]$	Average number of collisions in a virtual transmission time
$E[S]$	Average length of a successful transmission time
Idle_ $p_i$	$i^{th}$ idle period in a virtual transmission time
$f$	Traffic function
$M$	Number of immediate (1-hop away) neighbors of a node
$BT_{thre}$	Threshold for the back-off time
$l$	load
$\bar{m}$	Average message length
$n_{coll-avg}$	Average number of collisions
$n_{coll}$	Number of collisions
$N_c$	Average number of collisions in a virtual transmission time
$p$	Parameter to sample the back-off time from a geometric distribution
$p_{min}$	Minimum value for $p$

$q$	Packet length distribution
$t_{slot}$	Time slot
$t_v$	Virtual transmission time
$\alpha$	Busy rate
$\alpha_2$	Smoothing factor
$\tau$	Propagation delay
$\tau_{opt}$	Optimal channel access probability
$\rho_{max}$	Capacity (Maximum throughput)
$T_{receive}$	Time when a packet is successfully received
$T_{ready-to-send}$	Time when a packet is ready to send
$W_{init}$	Initial value for contention window

# CHAPTER 1

## INTRODUCTION

### 1.1 Background

With the extensive growth of a wide range of monitoring applications and phenomena such as public health care, target detection and tracking, environmental e just to name a few, wireless sensor networks have proved to be the only right solution for all these event-driven applications.

They exhibit a network of low-cost, low power nodes composed of elementary communication tools such as memory, processor and a radio or transceiver. These sensor nodes are usually scattered randomly in an area of interest to collect information about the event that is about to happen. The self-organized sensor nodes are expected to be autonomously operating ever after being deployed, leaving them in a situation of no external control [2].

A Medium Access Control (MAC) protocol organizes access of all nodes to a shared medium by deploying some scheduling algorithms to specify who and when can occupy the channel for the current transmission. The MAC design space for a WSN is limited by always considering an energy aware behaviour when enhancing any other traditional performance parameters such as latency, throughput and fairness. WSN-specific MAC protocols follow the basic design principles of IEEE 802.11 Carrier Sense Multiple Access with Collision Avoidance (CSMA/CA) protocol developed for wireless Local Area Networks (LANs). However due to the unique



properties and limitations of sensor nodes, much energy will be wasted from these reported sources: Idle listening, overhearing, collisions, traffic fluctuations and protocol overhead [2] , which will be later discussed in Chapter 2.

In this thesis we focus on providing an energy aware operation of the MAC protocol by reducing the collisions in Sensor MAC or S-MAC [Ye, Heidemann & Estrin. 2002] as the first MAC protocol for WSN which is designed on top of the IEEE 802.11 MAC protocol. IEEE 802.11 MAC protocol employs a Binary Exponential Back-off (BEB) algorithm to alleviate the probability of collision and adjusts the Contention Window (CW) for the next transmission based on an exponential behaviour which is discussed in depth in the next section. While the BEB mechanism in IEEE 802.11 changes the size of the CW exponentially the size of the CW in S-MAC is fixed. Many alternatives of MAC protocols for both WLANs and WSNs (mostly for WLANs) have been proposed introducing some novel techniques that can lead to a more proper size of the CW. In summary, while manipulating the back-off behaviour of IEEE 802.11 standard to have a more effective control on the usage of medium has been an active research area for years, there have been also remarkable attempts recently to address the problem of fixed CW in S-MAC in similar ways.

## **1.2 Problem Statement and Motivation**

Besides the primary task of creating the network infrastructure for establishing links between the randomly distributed nodes [3], MAC protocols also take control of the radio, the most power consuming component of a sensor node. In contrary to Bluetooth and mobile ad-hoc networks, as the closest peers to sensor networks, there



are larger number of nodes, a lower transmission power ( $\sim 0$  dBm) and also less radio range than that of Bluetooth or MANET [3].

Extensive researches such as those discussed later in the related literature brought in Section 2.8 in Chapter 2, show that the probability of collisions can be reduced to a great extent by a wise scheduling of the sensor node's transceiver, telling it when to access the channel so that it faces less collisions and thus less amount of delay. Collisions are only one of the sources of more produced latency and consequently a waste of energy in a WSN. Other sources reported in the literature are idle listening (as the most power consuming one), overhearing and control packet overhead which will be detailed later in Chapter 2.

A collision happens when two or more sensor nodes try to contend for the medium at the same time as there is only one common channel. It is obvious that when a collision happens extra time should be taken due to the retransmission of corrupted packets which leads to a smaller delivery ratio and overall throughput of the network. The probability of a collision will be reduced if each node contends for the medium in a proper time slot. IEEE 802.11 defines two methods for accessing the channel in a time slot; 1) Distributed Coordination Function (DCF) which is a contention-based one and 2) Point Coordination Function (PCF) which is a contention-free method. IEEE 802.11 MAC protocol DCF adopts two different schemes for sending and receiving data packets which are: 1) Two way handshaking also known as the basic CSMA/CA scheme and 2) Four way handshaking scheme also known as CSMA/CA with RTS/CTS. Both these schemes are also discussed in Chapter 2.