



UNIVERSITI PUTRA MALAYSIA

***MULTI-CRITERIA DIVISIBLE LOAD SCHEDULING IN BINARY TREE
NETWORK***

SHAMSOLLAH GHANBARI

IPM 2016 9



**MULTI-CRITERIA DIVISIBLE LOAD SCHEDULING IN BINARY TREE
NETWORK**

By

SHAMSOLLAH GHANBARI

**Thesis Submitted to the School of Graduate Studies, Universiti Putra Malaysia,
in Fulfilment of the Requirements for the Degree of Doctor of Philosophy**

January 2016

COPYRIGHT

All material contained within the thesis, including without limitation text, logos, icons, photographs, and all other artwork, is copyright material of Universiti Putra Malaysia unless otherwise stated. Use may be made of any material contained within the thesis for non-commercial purposes from the copyright holder. Commercial use of material may only be made with the express, prior, written permission of Universiti Putra Malaysia.

Copyright© Universiti Putra Malaysia



DEDICATION

To:

My father who is a root for me on the earth

My mother (her soul) who released me like a sapling from the dimness of the earth

My wife who turned me into a stalwart tree, and adorned me with the green leaves of
life

My children (Amir and Ilia) who magnified the beauty of my life a hundred times
like a fruit

All my professors and teachers who taught me the way of inflorescence and the
customs of true life



Abstract of thesis presented to the Senate of Universiti Putra Malaysia in fulfilment
of the requirement for the degree of Doctor of Philosophy

**MULTI-CRITERIA DIVISIBLE LOAD SCHEDULING IN BINARY TREE
NETWORK**

By

SHAMSOLLAH GHANBARI

January 2016

Chair : Professor Mohamed Othman, PhD
Institute : Mathematical Research

The Divisible Load Theory (DLT) is a paradigm in the area of parallel and distributed computing. Based on the divisible load theory, the computation and communication can be divided into some arbitrary independent parts, in which each part can be processed independently by a processor.

A class of the scheduling methods which is defined based on the DLT is called the Divisible Load Scheduling (DLS).

The traditional divisible load scheduling assumes that the processors report their true computation and communication rates, i.e., they do not cheat the algorithm. In the real applications, the processors may cheat the algorithm, which means, the processors might not report their true computation or communication rates. However, the problem that the processors may not report their true computation rates is called computation rate-cheating problem. The same definition can be considered for the communication rate-cheating problem. However, this problem was investigated by Thomas E. Carroll and D. Grosu in their research publications. The results of their research indicate that the computation rate-cheating reduces the performance of the divisible load scheduling.

This thesis focuses on the computation and communication rate-cheating problems aiming to reduce the effects of computation and communication rate-cheating on the performance of the divisible load scheduling model.

We adopt a multi-criteria approach to the problem. We propose three different multi-criteria based methods in order to improve the performance of the divisible load scheduling. The first method is a multi-objective divisible load scheduling method. The results show that this method is able to considerably improve the performance of the divisible load scheduling when the processors cheat their computation rates. The

experimental results indicate that the proposed method is able to reduce the finish time by approximately 66% in the best case. The limitation of the proposed multi-objective method is that this method slightly increases the start-up time. In order to reduce the limitation of the multi-objective method, a second method has been proposed, which is an Analytical Hierarchy Process (AHP) method. It is briefly called AHP-based method.

The experimental results show that the AHP-based method is able to improve the performance of divisible load scheduling. In addition, it has a lower start-up time comparing the multi-objective method.

In the third proposed method, it is assumed that both the communication and computation might not be reported at the true rates; hence, we have a new approach to the communication and computation rate-cheating problems. We propose a priority-based divisible load scheduling method for the first time. The results show that this method is able to allocate the optimal load when the processors cheat their computation and communication rates. The proposed priority-based divisible load scheduling method is a novel effort in the area of divisible load scheduling over the past two decades.

Abstrak tesis yang dikemukakan kepada Senat Universiti Putra Malaysia sebagai memenuhi keperluan untuk ijazah Doktor Falsafah

MULTI-KRITERIA BAGI PENJADUALAN BEBAN BERBAGI DALAM RANGKAIAN POKOK DEDUA

Oleh

SHAMSOLLAH GHANBARI

Januari 2016

Pengerusi : Professor Mohamed Othman, PhD
Institut : Penyelidikan Matematik

Teori Beban Berbagi (DBB) adalah paradigma dalam bidang pengkomputeraan selari dan teragih. Berdasarkan teori beban berbagi, pengiraan dan komunikasi boleh dibahagikan kepada sebarang bahagian yang bebas, dimana setiap bahagian boleh diproses secara berasingan oleh pemproses.

Terdapat satu kelas di dalam cara penjadualan yang ditakrifkan dalam konteks DBB yang dinamakan sebagai Penjadualan Beban Berbagi (PBB).

Penjadualan Beban Berbagi (PBB) tradisional menganggap bahawa pemproses-pemproses memberi laporan yang betul bagi kadar pengiraan dan kadar komunikasi, iaitu tidak berlaku penipuan algoritma. Di dalam aplikasi sebenar, pemproses mungkin menipu algoritma, bermakna pemproses-pemproses mungkin tidak melaporkan kadar pengiraan atau kadar komunikasi yang betul. Walau bagaimanapun, masalah dimana pemproses mungkin tidak memberikan laporan kadar pengiraan yang betul dinamakan sebagai masalah penipuan kadar pengiraan. Takrif yang sama diberikan kepada masalah penipuan kadar komunikasi. Namun, masalah ini telah dikaji oleh Thomas E. Carroll dan D. Grosu di dalam penerbitan penyelidikan mereka. Hasil penyelidikan mereka menunjukkan bahawa penipuan kadar pengiraan menyebabkan pengurangan prestasi bagi penjadualan beban berbagi.

Tesis ini memberi tumpuan kepada masalah penipuan kadar pengiraan dan komunikasi dalam mengurangkan kesan oleh penipuan kadar pengiraan dan kadar komunikasi ke atas prestasi model penjadualan beban berbagi.

Kami mempunyai pendekatan multi-kriteria kepada masalah tersebut. Kami mencadangkan tiga multi-kriteria berdasarkan kaedah dalam meningkatkan prestasi penjadualan beban berbagi.

Kaedah yang pertama adalah kaedah objektif berganda penjadualan beban berbagi. Keputusan menunjukkan bahawa kaedah ini mampu meningkatkan prestasi penjadualan beban berbagi apabila berlaku penipuan kadar pengiraan oleh pemproses. Keputusan eksperimen telah menunjukkan bahawa kaedah yang dicadangkan mampu untuk mengurangkan masa selesai kira-kira 66% dalam kes yang terbaik. Kaedah objektif berganda yang dicadangkan adalah terhad kerana kaedah ini meningkatkan sedikit masa permulaan.

Bagi mengurangkan had bagi kaedah objektif berganda, kaedah yang kedua telah dicadangkan, iaitu kaedah Proses Analisis Hirarki (PAH) yang dikenali sebagai kaedah PAH-asas.

Keputusan eksperimen menunjukkan bahawa kaedah ini juga dapat meningkatkan prestasi penjadualan beban berbagi. Di samping itu, kaedah ini telah dapat mengurangkan masa permulaan berbanding dengan kaedah objektif berganda.

Dalam kaedah ketiga yang dicadangkan, ia diandaikan bahawa kedua-dua komunikasi dan pengiraan mungkin tidak dilaporkan pada kadar benar; maka, kami memperkenalkan pendekatan baru kepada masalah penipuan kadar pengiraan dan komunikasi. Kami mencadangkan satu kaedah keutamaan-asas bagi kaedah penjadualan beban berbagi untuk kali pertama. Keputusan menunjukkan bahawa kaedah ini mampu memberi peruntukkan beban yang besar apabila berlaku penipuan kadar pengiraan dan komunikasi oleh pemproses. Kaedah keutamaan asas penjadualan beban berbagi yang dicadangkan adalah satu usaha murni di dalam bidang penjadualan beban berbagi sejak dua dekad yang lalu.

ACKNOWLEDGEMENTS

First of all, I would like to express my special gratitude to Allah who has supported me throughout my life. Specifically, He has supported me when I was surrounded with severe hardships and troubles during my Ph.D program.

I would like to express my sincere gratitude to my supervisor Professor Dr. Mohamed Othman who was supportive, helpful and patient during my doctoral studies. In fact, completing my Ph.D thesis would have been impossible without his cooperation, guidance and continuous assistance.

I would also like to thank Assoc. Prof. Dr. Mohd Rizam Abu Bakar and Assoc. Prof. Dr. Wah June Leong who were my supervisory committee members.

I would also like to express my heartfelt appreciation to my wife Maryam, and my son Amir Mohammad, for their love, supports, patience, time and efforts that they assigned me, my thesis and the Ph.D program. They helped me in all aspects from their psychological supports to their helpful attitude toward the thesis.

I also wish to express appreciation to my brothers Seifollah and Abdullah for their long distance, psychological and other helpful efforts and supports, which were truly vital during my study. Finally, I would like to express my deepest gratitude to my father for everything.

I certify that a Thesis Examination Committee has met on 12 January 2016 to conduct the final examination of Shamsollah Ghanbari on his thesis entitled "Multi-Criteria Divisible Load Scheduling in Binary Tree Network" in accordance with the Universities and University Colleges Act 1971 and the Constitution of the Universiti Putra Malaysia [P.U.(A) 106] 15 March 1998. The Committee recommends that the student be awarded the Doctor of Philosophy.

Members of the Thesis Examination Committee were as follows:

Adem Kilicman, PhD

Professor
Faculty of Science
Universiti Putra Malaysia
(Chairman)

Zanariah binti Abdul Majid, PhD

Professor
Faculty of Science
Universiti Putra Malaysia
(Internal Examiner)

Shaharuddin Salleh, PhD

Professor
Universiti Teknologi Malaysia
Malaysia
(External Examiner)

Debasish Ghose, PhD

Professor
Indian Institute of Science
India
(External Examiner)



ZULKARNAIN ZAINAL, PhD

Professor and Deputy Dean
School of Graduate Studies
Universiti Putra Malaysia

Date: 16 February 2016

This thesis was submitted to the Senate of Universiti Putra Malaysia and has been accepted as fulfilment of the requirement for the degree of Doctor of Philosophy. The members of the Supervisory Committee were as follows:

Mohamed Othman, PhD

Professor

Faculty of Computer Science and Information Technology

(Institute for Mathematical Research)

Universiti Putra Malaysia

(Chairperson)

Mohd Rizam Abu Bakar, PhD

Associate Professor

Faculty of Science

Universiti Putra Malaysia

(Member)

Wah June Leong, PhD

Associate Professor

Faculty of Science

Universiti Putra Malaysia

(Member)

BUJANG KIM HUAT, PhD

Professor and Dean

School of Graduate Studies

Universiti Putra Malaysia

Date:

Declaration by graduate student

I hereby confirm that:

- this thesis is my original work
- quotations, illustrations and citations have been duly referenced
- the thesis has not been submitted previously or concurrently for any other degree at any institutions
- intellectual property from the thesis and copyright of thesis are fully-owned by Universiti Putra Malaysia, as according to the Universiti Putra Malaysia (Research) Rules 2012;
- written permission must be owned from supervisor and deputy vice –chancellor (Research and innovation) before thesis is published (in the form of written, printed or in electronic form) including books, journals, modules, proceedings, popular writings, seminar papers, manuscripts, posters, reports, lecture notes, learning modules or any other materials as stated in the Universiti Putra Malaysia (Research) Rules 2012;
- there is no plagiarism or data falsification/fabrication in the thesis, and scholarly integrity is upheld as according to the Universiti Putra Malaysia (Graduate Studies) Rules 2003 (Revision 2012-2013) and the Universiti Putra Malaysia (Research) Rules 2012. The thesis has undergone plagiarism detection software

Signature: _____ Date: _____

Name and Matric No: Shamsollah Ghanbari, GS31795

Declaration by Members of Supervisory Committee

This is to confirm that:

- the research conducted and the writing of this thesis was under our supervision;
- supervision responsibilities as stated in the Universiti Putra Malaysia (Graduate Studies) Rules 2003 (Revision 2012-2013) were adhered to.

Signature: _____
Name of
Chairman of
Supervisory
Committee: Professor Mohamed Othman

Signature: _____
Name of
Member of
Supervisory
Committee: Assoc. Prof Mohd Rizam Abu Bakar

Signature: _____
Name of
Member of
Supervisory
Committee: Assoc. Prof. Wah June Leong

TABLE OF CONTENTS

	Page
ABSTRACT	i
ABSTRAK	iii
ACKNOWLEDGEMENTS	v
APPROVAL	vi
DECLARATION	viii
LIST OF TABLES	xiii
LIST OF FIGURES	xv
LIST OF ABBREVIATIONS	xviii
CHAPTER	
1. INTRODUCTION	1
1.1. Background	1
1.2. Problem Statement	2
1.3. Research Objectives	2
1.4. Research Scope	3
1.5. Research Significance	3
1.6. Research Contributions	4
1.7. Thesis Organization	5
2. LITERATURE REVIEW	6
2.1. Introduction	6
2.2. Divisible Load Scheduling Model	6
2.2.1. Single Level Tree Network	7
2.2.2. Multi-level Tree Network	8
2.2.3. Other Topologies	8
2.3. Basic Strategies for Solving the DLS Problem	9
2.3.1. Markov Chain Model	9
2.3.2. Linear Programming Model	10
2.3.3. Non-linear Model	10
2.4. Advanced Strategies for Solving the DLS Problem	10
2.4.1. Multi-installment Processing	11
2.4.2. Adaptive Strategy	12
2.4.3. Multiple-source Allocation Model	12
2.5. Multi-criteria Approach to the DLS	13
2.5.1. Multi-objective Decision-making	13
2.5.2. Multi-attribute Decision-making Problem	13
2.6. Applications of Divisible Load Theory	13
2.7. Critical Analysis of the Research in the Area of DLS	14
2.7.1. Typology of the Existing Research	14
2.7.2. Algorithms	17
2.7.3. Related Works	19
2.8. Summary	20

3.	METHODOLOGY	21
3.1.	Introduction	21
3.2.	Preliminaries	21
3.2.1.	Definition of Parameters	21
3.2.2.	Performance Metrics	22
3.2.3.	Benchmarking	23
3.3.	General Description of the Proposed Methods	23
3.3.1.	Multi-objective Method	26
3.3.2.	AHP-based Method	30
3.3.3.	Priority-based Method	31
3.4.	Performance Evaluation	32
3.4.1.	Experimental Setup	32
3.4.2.	Analysis of the Results	32
3.5.	Summary	32
4.	MULTI-OBJECTIVE DIVISIBLE LOAD SCHEDULING IN A BINARY TREE NETWORK	34
4.1.	Introduction	34
4.2.	Assumptions	34
4.3.	Notations and Definitions	34
4.4.	Proposed Method	35
4.4.1.	Initializing	36
4.4.2.	Computation-based Probing	36
4.4.3.	Approximation	38
4.4.4.	Optimization	39
4.4.5.	Load Allocation	41
4.5.	Performance Evaluation	41
4.5.1.	Experimental Setup	41
4.5.2.	Simulation Method	41
4.5.3.	Simulation Results	44
4.6.	Discussion	47
4.6.1.	Analysis of the Results	47
4.6.2.	Effects of the Rate-cheating on the Finish Time	53
4.6.3.	Effects of the Rate-cheating on the Speed-up	55
4.6.4.	Limitation of the Proposed Method	57
4.7.	Summary	58
5.	ANALYTICAL HIERARCHY PROCESS METHOD FOR THE DIVISIBLE LOAD SCHEDULING	60
5.1.	Introduction	60
5.2.	Analytical Hierarchy Process	60
5.3.	Assumptions	61
5.4.	Notations and Definitions	61
5.5.	Proposed Method	62
5.5.1.	Computation-based Probing	63
5.5.2.	Approximating the Level of Cheating	65
5.5.3.	Estimation of the Actual Computation Rate	67
5.5.4.	Load Allocation	68

5.6.	Performance Evaluation	69
5.7.	Simulation Results and Discussion	70
5.7.1.	Proposed Method with Limited Cheat-processors	70
5.7.2.	Proposed Method with Unlimited Cheat-processors	78
5.8.	Summary	81
6.	PRIORITY-BASED DIVISIBLE LOAD SCHEDULING IN A SINGLE-LEVEL TREE NETWORK	83
6.1.	Introduction	83
6.2.	Assumptions	83
6.3.	Notations and Definitions	83
6.4.	Problem Description	84
6.5.	Framework	85
6.6.	Proposed Method	87
6.6.1.	Communication-based Probing	88
6.6.2.	Computation-based Probing	89
6.6.3.	Calculation of the Priority of Processors	90
6.6.4.	Load Allocation	92
6.7.	Experimental Results and Discussion	92
6.8.	Summary	99
7.	CONCLUSION AND FUTURE WORKS	100
7.1.	Conclusion	100
7.2.	Future Works	101
	REFERENCES	102
	BIODATA OF STUDENT	109
	LIST OF PUBLICATIONS	110

LIST OF TABLES

Table	Page
2.1. Important application topics of the DLS	14
2.2. A list of important topics of study in the area of DLS	15
2.3. Typology of research in the area of DLS	17
2.4. Comparing the algorithms applied for the DLS	18
3.1. Comparing the results of re-implementation and Carroll's results (balanced binary tree in fast rate)	24
3.2. Comparing the results of re-implementation and Carroll's results (balanced binary tree in slow rate)	24
3.3. Comparing the results of re-implementation and Carroll's results (left unbalanced binary tree in slow rate)	27
3.4. Comparing the results of re-implementation and Carroll's results (right unbalanced binary tree in slow rate)	27
4.1. Definitions and notations	34
4.2. A sample of population produced by the computation-based probing.	37
4.3. Information about the experimental	41
4.4. True computation rate for 15 processors	42
4.5. The reported, actual and estimated computation rates in fast cheat	43
4.6. The reported, actual and estimated computation rates in slow cheat.	43
4.7. Comparing the results of proposed method and Carroll's method (balanced binary tree in fast rate)	46
4.8. Comparing the results of proposed method and Carroll's method (balanced binary tree in slow rate)	48
4.9. Comparing the results of proposed method and Carroll's method (right unbalanced binary tree in slow rate)	50
4.10. Comparing the results of proposed method and Carroll's method (left unbalanced binary tree in slow rate)	52
5.1. Random index vs number of rows (columns) of matrix	61
5.2. Definitions and notations	61
5.3. Sample of population produced by the computation-based probing.	64
5.4. A rule-based for estimating actual computation rate	68
5.5. The reported, actual and estimated computation rates in fast cheat	70
5.6. The reported, actual and estimated computation rates in slow cheat	70
5.7. Comparing the total finish time of the AHP-based method with the multi-objective method and Carroll's results in fast rate for balanced binary tree	71
5.8. Comparing the total finish time of AHP-based method with the multi- objective method and Carroll's results in slow rate for balanced binary tree	71
5.9. Comparing the total finish time of proposed method with the multi- objective method and Carroll's results in slow rate for left unbalanced binary tree	75
5.10. Comparing the total finish time of proposed method and Carroll's results in slow rate for right unbalanced binary tree	77
6.1. Definitions and notations	84
6.2. Sample of population produced by the communication-based probing	88
6.3. Sample of population produced by the computation-based probing.	89

6.4.	The population produced by communication-based probing	95
6.5.	Comparison matrix of the processors in the first communication-based probing	96
6.6.	Comparison matrix of the processors in the second communication-based probing.	96
6.7.	Comparison matrix of the processors in the third communication-based probing	96
6.8.	Comparison matrix of the processors in the fourth communication-based probing	96
6.9.	The population produced by computation-based probing	97



LIST OF FIGURES

Figure	Page
1.1. The scope of research	4
2.1. A typical form of load distribution on the single level tree network.	7
2.2. Gantt chart-like timing diagram for the DLS in a single-level tree network	7
2.3. Gantt chart-like timing diagram for the DLS in a multi-level tree network	9
2.4. Morkov model for the DLS with a single-level tree network topology based on Moges and Robertazzi (2006)	9
2.5. Gantt chart-like diagram for the multi-installment DLS	11
2.6. Typology of the research in the area of the DLS (the number of citation)	16
2.7. Typology of the research in the area of the DLS (the number of papers)	16
3.1. Comparing the results of re-implementation and Carroll's results (the total finish time of the balanced binary tree in fast rate)	24
3.2. Comparing the results of re-implementation and Carroll's results (the utility of the balanced binary tree in fast rate)	25
3.3. Comparing the results of re-implementation and Carroll's results (the payment of the balanced binary tree in fast rate)	25
3.4. Comparing the results of re-implementation and Carroll's results (the total fi time of the balanced binary tree in slow rate)	26
3.5. Comparing the results of re-implementation and Carroll's results (the payment of the balanced binary tree in slow rate)	26
3.6. Comparing the results of re-implementation and Carroll's results (the utility of the balanced binary tree in slow rate)	27
3.7. Comparing the results of re-implementation and Carroll's results (the total finish time of the left unbalanced tree in slow rate)	28
3.8. Comparing the results of re-implementation and Carroll's results (the payment of the left unbalanced tree in slow rate)	28
3.9. Comparing the results of re-implementation and Carroll's results (the utility right unbalanced tree in slow rate)	29
3.10. Comparing the results of re-implementation and Carroll's results (the total finish time of the right unbalanced tree in slow rate)	29
3.11. Comparing the results of re-implementation and Carroll's results (the payment of the right unbalanced tree in slow rate)	30
3.12. Comparing the results of re-implementation and Carroll's results (the utility of the left unbalanced tree in slow rate)	30
3.13. A general framework for the research	31
4.1. A general framework for the proposed multi-objective method	35
4.2. Balanced binary tree network with 15 processors	42
4.3. Right unbalanced binary tree network with 15 processors	42
4.4. Left unbalanced binary tree network with 15 processors	42
4.5. The sub-trees of the third level the of the balanced binary tree network.	43
4.6. The sub-trees of the second level of the balanced binary tree network.	44

4.7.	The sub-tree of the first balanced binary tree network.	44
4.8.	Comparing the finish time of the proposed method and Carroll's method (balanced binary tree in fast rate)	45
4.9.	Comparing the payment of the proposed method and Carroll's method (balanced binary tree in fast rate)	45
4.10.	Comparing the utility of the proposed method and Carroll's method (balanced binary tree in fast rate)	46
4.11.	Comparing the finish time of the proposed method and Carroll's method (balanced binary tree in slow rate)	47
4.12.	Comparing the payment of the proposed method and Carroll's method (balanced binary tree in slow rate)	47
4.13.	Comparing the utility of the proposed method and Carroll's method (balanced binary tree in slow rate)	48
4.14.	Comparing the finish time of the proposed method and Carroll's method (right unbalanced binary tree in slow rate)	49
4.15.	Comparing the payment of the proposed method and Carroll's method (right unbalanced binary tree in slow rate)	49
4.16.	Comparing the utility of the proposed method and Carroll's method (right unbalanced binary tree in slow rate)	50
4.17.	Comparing the finish time of the proposed method and Carroll's method (left unbalanced binary tree in slow rate)	51
4.18.	Comparing the payment of the proposed method and Carroll's method (left unbalanced binary tree in slow rate)	51
4.19.	Comparing the utility of the proposed method and Carroll's method (left unbalanced binary tree in slow rate)	52
4.20.	The effects of computation rate-cheating on the makespan (p_1 is in the high cheat mode)	53
4.21.	The effects of computation rate-cheating on the makespan (p_1 is in the low cheat mode)	54
4.22.	The effect of cheated rate on the idle time in fast rate	54
4.23.	The effect of cheated rate on the idle time in slow rate	55
4.24.	Comparing the speed-up of the proposed method and Carroll's method (balanced binary tree in fast rate)	56
4.25.	Comparing the speed-up of the proposed method and Carroll's method (balanced binary tree in slow rate)	56
4.26.	Comparing the speed-up of the proposed method and Carroll's method (left unbalanced binary tree)	57
4.27.	Comparing the speed-up of the proposed method and Carroll's method (right unbalanced binary tree)	57
5.1.	A general framework for the proposed AHP-based method	63
5.2.	Empirical distribution function for the AHP-based method in fast rate	69
5.3.	Comparing the finish time of the proposed methods with Carroll's method (balanced binary tree in fast rate)	71
5.4.	Comparing the payment of the proposed methods with Carroll's method (balanced binary tree in fast rate)	72
5.5.	Comparing the utility of the proposed methods with Carroll's method (balanced binary tree in fast rate)	72
5.6.	Comparing the finish time of the proposed methods with Carroll's method (balanced binary tree in slow rate)	73

5.7.	Comparing the payment of the proposed methods with Carroll's method (balanced binary tree in slow rate)	73
5.8.	Comparing utility time of the proposed methods with Carroll's method (balanced binary tree in slow rate)	74
5.9.	Comparing the fin time of the proposed methods with Carroll's method (right unbalanced binary tree in slow rate.)	74
5.10.	Comparing the payment of the AHP-based methods with Carroll's method (right unbalanced binary tree in slow rate)	75
5.11.	Comparing the utility of the proposed methods with Carroll's method (right unbalanced binary tree in slow rate)	76
5.12.	Comparing the finish time of the proposed methods with Carroll's method (left unbalanced binary tree in slow rate)	76
5.13.	Comparing the payment of the proposed methods with Carroll's method (left unbalanced binary tree in slow rate)	77
5.14.	Comparing the utility of the proposed methods with Carroll's method (left unbalanced binary tree in slow rate)	78
5.15.	Comparing the finish time plus the start-up time in two proposed methods and Carroll's method in case 1	79
5.16.	Comparing the finish time plus the start-up time in two proposed methods and Carroll's method in case 2	79
5.17.	Comparing the finish time plus the start-up time in two proposed methods and Carroll's method in case 6 and 7	80
5.18.	Comparing the finish time plus the start-up time in two proposed methods and Carroll's method in case 8	80
5.19.	Comparing the finish time plus the start-up time in two proposed methods and Carroll's method in case 3	81
5.20.	Comparing the finish time plus the start-up time in two proposed methods and Carroll's method in cases 4 and 5	81
6.1.	A hierarchical framework for the proposed priority-based method	86
6.2.	A descriptive framework for the proposed priority-based method	87
6.3.	A single-level tree network with five processors	93
6.4.	The priority of processors in three probing	94
6.5.	Processing time in different communication-based probing processes	95
6.6.	The priority of processors based on communication criteria	97
6.7.	The priority of processors based on computation criteria	98
6.8.	Evaluation of the priority-based method in 24 different cases	99

LIST OF ABBREVIATIONS

AHP	Analytical Hierarchy Process
CI	Consistency Index
CR	Consistency Rate
DLS	Divisible Load Scheduling
DLT	Divisible Load Theory
EDF	Empirical Distribution Function
GA	Genetic Algorithm
IDLT	Iterative Divisible Load Theory
MADM	Multi-Attribute Decision Making
MCDM	Multi-Criteria Decision Making
NFT	Normal Form Table
NP	Non Polynomial
PCD	Probing and Continuous Distribution Strategy
PDD	Probing and Delayed Distribution Strategy
PSD	Probing and Selective Distribution Strategy
RI	Random Index
RUMR	Robust-UMR
UMR	Uniform Multi-Round

CHAPTER 1

INTRODUCTION

1.1 Background

The problem of how to effectively distribute the jobs among resources to improve the performance so that some jobs do not suffer unbounded delays is called job scheduling. It is very important to assign appropriate resources to the jobs. Through a good scheduling algorithm, the system can perform better and applications can avoid unnecessary delays.

A class of scheduling techniques that is suitable for using in the area of parallel and distributed computing with big data is called Divisible Load Scheduling (DLS). The first article concerning the Divisible Load Theory (DLT) was proposed by Yuan-Chieh and Robertazzi (1988). In the same time Agrawal and Jagadish (1988) proposed another divisible load scheduling algorithm in a separate paper. The DLT is based on the fact that the load can be divided into various parts, in which each part can be executed by an independent processor. Over the past two decades, the DLT has found a wide variety of applications in the area of parallel processing. Moreover, the DLT has been applied to a wide variety of interconnection topologies. Basically, it was applied to the bus and single level tree network topologies. Bataineh et al. (1994) used the DLT for a multi level tree and daisy chain. After that Blazewicz and Drozdowski (1995) developed the DLT for hypercubes. Later, the DLT was applied to the two-dimensional, three-dimensional and k-dimensional meshes by Blazewicz et al. (1999a), Drozdowski and Glazek (1999) and Li (2003) respectively. It has also been applied in homogeneous and heterogeneous platforms by Blazewicz et al. (1999b) and Beaumont et al. (2003) respectively. More recently, the DLT was used for scheduling in grid and cloud environments by Robertazzi (2007) and Suresh et al. (2015) respectively.

In general, the divisible load scheduling assumes that the initial amount V of the load is held by the originator denoted by p_0 . A common assumption is that the originator does not do any computation. It only distributes the load into parts $\alpha_1, \alpha_2, \dots, \alpha_m$ to be processed on the worker processors denoted by p_1, p_2, \dots, p_m . According to Sohn and Robertazzi (1993), the condition for the optimal solution is that, all of the processors stop processing at the same time; otherwise, the load could be transferred from busy to idle processors to improve the solution time. Assume that w_0, w_1, \dots, w_m are the inverse computing speeds (computation rates) of the processors. It is also assumed that z_0, z_1, \dots, z_m are the inverse transmission speeds (communication rates) of the processors. In this situation we assume that the transmission commences simultaneously on all links, and that computation follows the load reception on each processor.

The traditional divisible load scheduling assumes that the processors report their true computation and communication rates, i.e., they do not cheat the algorithm. In the real applications, the processors may cheat the algorithm. This means that the processors may not report their true computation or communication rates.

This issue was investigated by Carroll and Grosu (2008, 2012) in their research publications. The results of their research indicate that the computation rate-cheating, reduces the performance of the divisible load scheduling. In fact, the divisible load scheduling model only obtains its optimal performance if the processors report their true computation rates.

This thesis focuses on the computation and communication rate-cheating problems. We have a multi-criteria perspective to the problem. The goal is to improve the performance of the divisible load scheduling, where the processors cheat the algorithm.

1.2 Problem Statement

Based on the traditional divisible load theory, it is assumed that, the processors do not cheat the algorithm. This means that the traditional divisible load theory assumes that, the processors report their true computation and communication rates to the originator.

The divisible load scheduling has also been examined on the processors that cheat the algorithm. The main idea of the computation and communication rate-cheating problems refer to the misreporting and time varying problems, which were investigated in respect of the divisible load scheduling by Jeeho and Robertazzi (1998). Subsequently, Carroll and Grosu (2008, 2012) focused on the application case of misreporting in the divisible load scheduling. They proposed a strategy-proof mechanism for the divisible load scheduling under the bus and single-level network topologies. A few years later, they investigated the effects of computation rate-cheating on the multi-level tree network topology. Finally, they proposed an incentive-based mechanism for the divisible load scheduling under the multi-level tree network topology.

However, the computation rate-cheating problem may occur, if the processors execute their fraction of loads with different rates. Suppose that, the originator allocates $\alpha = (\alpha_0, \alpha_1, \dots, \alpha_m)$ fraction of load to the processors. This allocation is based on the assumption that, the actual computation rate of p_j is equal to w_j . It is also assumed that the actual communication rate of p_j is equal to z_j . In fact, the originator (p_0) learns the actual computation rate of the worker processors once they complete execution of their fraction of the load. According to Carroll and Grosu (2008, 2012), if the processors cheat the algorithm the divisible load scheduling model will fail to achieve its optimal performance. As a result, the total finish time will be increased on this occasion.

This thesis focuses on the computation and communication rate-cheating problems with the aim of reducing the effects of rate-cheating on the performance of the divisible load scheduling.

1.3 Research Objectives

The goal of this research is to present scheduling based on the divisible load theory in order to obtain the best possible response time when the processors cheat the system.

The objectives of this research are to

- propose a multi-objective divisible load scheduling method in a binary tree network.
- propose an analytical hierarchy process divisible load scheduling method in a binary tree network.
- propose a priority-based divisible load scheduling in a tree network.

The mentioned objectives have the following effects on the existing research in the area of the divisible load scheduling.

- This research improves the performance of the divisible load scheduling when the processors cheat the algorithm.
- It secures the divisible load scheduling model against the computation rate-cheating problem.
- It also reduces the limitations of the previous works which was proposed by Carroll and Grosu (2008, 2012).
- Furthermore, this research reduces the effects of communication and computation rate-cheating on the finish time in a divisible load scheduling model.

1.4 Research Scope

This study is concerned with scheduling in the area of parallel and distributed computing. It focuses on the divisible load scheduling in the binary tree network topology. It is assumed that the worker processors may cheat the algorithm. The thesis also concerns the multi-criteria problem solving as a technique for improving the performance of the system under the computation and communication rate-cheating problems. It mainly uses the multi-objective and multi-attribute methods. We also focus on the analytical hierarchy process as a multi-attribute technique for handling the computation and communication rate-cheating problems. The scope of research is shown in Fig. 1.1. As the figure shows, by using probing process, the divisible load model is able to produce some information about the behaviour of the processors. The gathered information helps the algorithm to anticipate the actual computation and communication rates of the processors. For this purpose we use the multi-criteria decision making models.

1.5 Research Significance

Generally, this thesis focuses on the divisible load scheduling in the multi-level tree and single-level tree network topologies. It discusses the divisible load scheduling

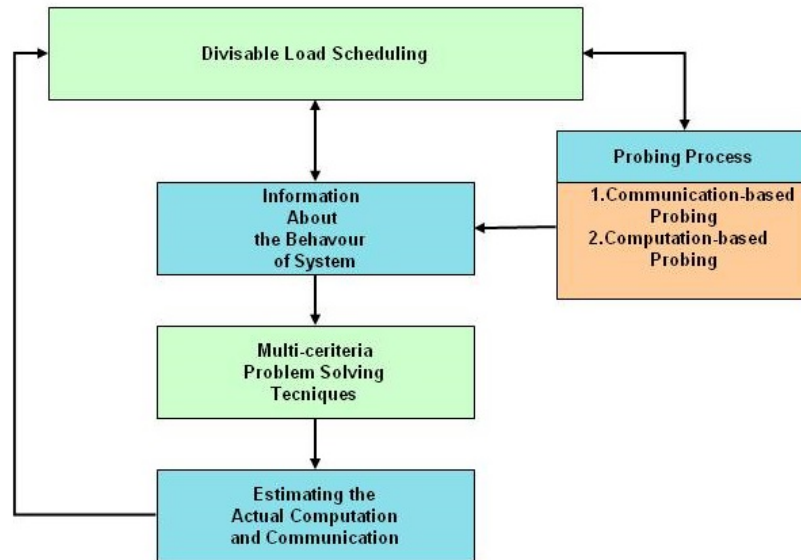


Figure 1.1: The scope of research.

under the communication and computation rate-cheating problem. The computation and communication rate-cheating problems are the important issues in the area of distributed and parallel computing. According to the recent research the rate-cheating decreases the performance of the divisible load scheduling model. This research improves the performance of divisible load scheduling model when the processors cheat their communication and computation rates.

1.6 Research Contributions

The following list demonstrates the contributions of the research:

- This thesis proposed a multi-objective divisible load scheduling method.
 - The proposed multi-objective method enhanced the performance of the divisible load scheduling when the processors cheat the algorithm.
 - It improved the total finish time, payment and utility.
 - It also reduced the finish time by approximately 66% in the best case.
 - Moreover, we proposed a new method of computation-based probing which is useful for predicting the behaviour of divisible load scheduling.
- This research formulated the divisible load scheduling model as a multi-attribute decision-making problem for the first time.
 - We formulated the divisible load scheduling as an analytical hierarchy process problem.

- It improved the total finish time, payment and utility almost the same as the multi-objective method.
 - It also considerably reduced the start-up time.
- This research proposed a priority-based method for the divisible load scheduling for the first time.
 - The priority-based approach is a novel effort in the area of divisible load scheduling over the past two decades.
 - We proposed a new method of communication-based probing which is useful for predicting the behaviour of divisible load scheduling.
 - We also found a relationship between the comparison matrices and the processors. Since, the divisible load theory assumes that, the processing must be executed by the independent processors, the relationship between the comparison matrices and the processors helps us to predict the effects of each processor on the other processors.

1.7 Thesis Organization

The other chapters of this thesis are organized as follows. Chapter 2 presents a review of the literature concerning the divisible load theory. Chapter 3 generally describes the methodology used in this thesis. It also briefly explains the proposed methods. Three different methods are proposed in this thesis. The first method is a multi-objective divisible load scheduling which is explained in chapter 4. The second proposed method is an AHP-based divisible load scheduling method which is explained in chapter 5. The third proposed method is a priority-based divisible load scheduling which is explained in chapter 6. Lastly, the conclusion and future works are presented in chapter 7.



© COPYRIGHT UPM

REFERENCES

- Abdullah, M., Othman, M., Ibrahim, H. and Subramaniam, S. 2007. An integrated approach for scheduling divisible load on large scale data grids. *Lecture Notes in Computer Science* 4705 (1): 748–757.
- Abdullah, M., Othman, M., Ibrahim, H. and Subramaniam, S. 2010. Optimal workload allocation model for scheduling divisible data grid applications. *Future Generation Computer Systems* 26 (7): 971–978.
- Agrawal, R. and Jagadish, H. 1988. Partitioning techniques for large-grained parallelism. *IEEE Transactions on Computers* 37 (12): 1627–1634.
- Bataineh, S., Hsiung, T. Y. and Robertazzi, T. G. 1994. Closed form solutions for bus and tree networks of processors load sharing a divisible job. *IEEE Transactions on Computers* 43 (10): 1184–1196.
- Beaumont, O., Legrand, A. and Robert, Y. 2003. Scheduling divisible workloads on heterogeneous platforms. *Parallel Computing* 29 (9): 1121–1152.
- Berlinska, J. and Drozdowski, M. 2010. Heuristics for multi-round divisible loads scheduling with limited memory. *Parallel Computing* 36 (4): 199–211.
- Blazewicz, J. and Drozdowski, M. 1995. Scheduling divisible jobs on hypercubes. *Parallel Computing* 21 (12): 1945–1956.
- Blazewicz, J., Drozdowski, M., Guinand, F. and Trystram, D. 1999a. Scheduling a divisible task in a 2-dimensional Mesh. *Discrete Applied Mathematics* 94 (1-3): 35–50.
- Blazewicz, J., Drozdowski, M. and Markiewicz, M. 1999b. Divisible task scheduling concept and verification. *Parallel Computing* 25 (1): 87–98.
- Carroll, T. E. and Grosu, D. 2007. A strategyproof mechanism for scheduling divisible loads in linear networks. *IEEE International Parallel and Distributed Processing Symposium* 1–9.
- Carroll, T. E. and Grosu, D. 2008. Strategyproof mechanisms for scheduling divisible loads in bus-networked distributed systems. *IEEE Transactions on Parallel and Distributed Systems* 19 (8): 1124–1135.
- Carroll, T. E. and Grosu, D. 2012. An incentive-based distributed mechanism for scheduling divisible loads in tree-networks. *Parallel and Distributed Computing* 72: 389–401.
- Chan, S., Veeravalli, B. and Ghose, D. 2001. Large matrix-vector products on distributed bus networks with communication delays using the divisible load paradigm: performance analysis and simulation. *Mathematics and Computers in Simulation* 58 (1): 71–92.
- Chang, Y. K., Wu, J. H., Chen, C. Y. and Chu, C. P. 2007. Improved methods for divisible load distribution on k-dimensional meshes using multi-installment. *IEEE Transactions on Parallel and Distributed Systems* 18 (11): 1618–1629.

- Chen, P. S., Chu, P. and Lin, M. 2002. On vargas's proof of consistency test for 3x3 comparison matrices in AHP. *Journal of the Operations Research Society of Japan* 45 (3): 233–242.
- Chen, S. 2013. Adaptive indexed divisible load theory for wireless sensor network workload allocation. *International Journal of Distributed Sensor Networks*. doi: <http://dx.doi.org/10.1155/2013/484796> .
- Cheng, Y. and Robertazzi, T. G. 1990. Distributed computation for a tree network with communication delays. *IEEE Transactions on Aerospace and Electronic Systems* 26 (3): 511–516.
- Dantong, Y. and Robertazzi, T. G. 2003. Divisible load scheduling for grid computing. *Proceeding of International Conference on Parallel and Distributed Computing and Systems* 1: 1–6.
- Deluka-Tibljias, A., Karleuvs, B. and Dragivcevic, N. 2013. Review of multicriteria-analysis methods application in decision making about transport infrastructure. *Gradjevinar* 65 (7): 619–631.
- Drozdowski, M. 2009. *Scheduling for parallel processing. First edition*. Springer London.
- Drozdowski, M. and Glazek, W. 1999. Scheduling divisible loads in a three-dimensional mesh of processors. *Parallel Computing* 25 (4): 381–404.
- Drozdowski, M. and Lawenda, M. 2006. Multi-installment divisible load processing in heterogeneous systems with limited memory. *Parallel Processing and Applied Mathematics* 3911: 847–854.
- Drozdowski, M. and Wolniewicz, P. 2006. Optimum divisible load scheduling on heterogeneous stars with limited memory. *European Journal of Operational Research* 172 (2): 545–559.
- Eric, J. L. and Jean, S. 1982. Assessing a set of additive utility functions for multi-criteria decision-making. *European Journal of Operational Research* 10 (2): 151–164.
- Farahani, R. Z., SteadieSeifi, M. and Asgari, N. 2010. Multiple-criteria facility location problems: A survey. *Applied Mathematical Modeling* 34 (7): 1689–1709.
- Gerald, E. 1984. An overview of techniques for solving multi-objective mathematical programs. *Management Science* 30 (11): 1268–1282.
- Ghose, D. 2002. A feedback strategy for load allocation in workstation clusters with unknown network resource capabilities using the DLT paradigm. *Proceedings of the International Conference on Parallel and Distributed Processing Techniques and Applications* 1: 425–428.
- Ghose, D., Kim, H. J. and Kim, T. H. 2005. Adaptive divisible load scheduling strategies for workstation clusters with unknown network resources. *IEEE Transactions on Parallel and Distributed Systems* 16 (10): 897–907.

- Hillier, F. S., Lieberman, G. J., Nagerald, B. and Basu, P. 2012. *Introduction to operations research. 9th edition*. Tata McGraw-Hill Education.
- Hung, J. T. and Robertazzi, T. G. 2004. Divisible load cut through switching in sequential tree networks. *IEEE Transactions on Aerospace and Electronic Systems* 40 (3): 968–982.
- Hung, J. T. and Robertazzi, T. G. 2008. Scheduling nonlinear computational loads. *IEEE Transactions on Aerospace and Electronic Systems* 44 (3): 1169–1182.
- Jeeho, S. and Robertazzi, T. G. 1998. Optimal time-varying load sharing for divisible loads. *IEEE Transactions on Aerospace and Electronic Systems* 34 (3): 907–923.
- Jiang-Xia, N., Deng-Feng, L. and Mao-Jun, Z. 2010. A lexicographic method for matrix games with payoffs of triangular intuitionistic fuzzy numbers. *International Journal of Computational Intelligence Systems* 3 (3): 280–289.
- Kim, H. J. 2003. A novel optimal load distribution algorithm for divisible loads. *Cluster Computing* 6 (1): 41–46.
- Kim, S. and Weissman, J. B. 2004. A genetic algorithm based approach for scheduling decomposable data grid applications. *International Conference on Parallel Processing* 406–413.
- Ko, K. and Robertazzi, T. G. 2002. Scheduling in an environment of multiple job submissions. *Proceedings of the International Conference on Information Sciences and Systems* .
- Ko, K. and Robertazzi, T. G. 2004. Equal allocation scheduling for data intensive applications. *IEEE Transactions on Aerospace and Electronic Systems* 40 (2): 695–705.
- Ko, K. and Robertazzi, T. G. 2008. Signature search time evaluation in flat file databases. *IEEE Transactions on Aerospace and Electronic Systems* 44 (2): 493–502.
- Kuřakowski, K. 2015. Notes on order preservation and consistency in AHP. *European Journal of Operational Research* 245 (1): 333–337.
- Li, K. 2003. Improved methods for divisible load distribution on k-dimensional meshes using pipelined communications. *IEEE Transactions on Parallel and Distributed Systems* 14 (12): 1250–1261.
- Li, K. 2011. New divisible load distribution methods using pipelined communication techniques on tree and pyramid networks. *IEEE Transactions on Aerospace and Electronic Systems* 47 (2): 806–819.
- Li, P., Veeravalli, B. and Kassim, A. A. 2005. Design and implementation of parallel video encoding strategies using divisible load analysis. *IEEE Transactions on Circuits and Systems for Video Technology* 15 (9): 1098–1112.
- Li, X. and Veeravalli, B. 2010. PPDD: scheduling multi-site divisible loads in single-level tree networks. *Cluster Computing* 13 (1): 31–46.

- Li, X., Veeravalli, B. and Chung, C. 2000. Divisible load scheduling on single-level tree networks with buffer constraints. *IEEE Transactions on Aerospace and Electronic Systems* 36 (4): 1298–1308.
- Li, X., Veeravalli, B. and Ko, C. C. 2001. Divisible load scheduling on a hypercube cluster with finite-size buffers and granularity constraints. *The First IEEE/ACM International Symposium on Cluster Computing and the Grid* 660–667.
- Lin, X., Lu, Y., Deogun, J. and Goddard, S. 2007. Enhanced real-time divisible load scheduling with different processor available times. *Proceedings of the International Conference of High Performance Computing* 308–319.
- Mamat, A., Lu, Y., Deogun, J. and Goddard, S. 2012. Scheduling real-time divisible loads with advance reservations. *Real-time Systems* 48 (3): 264–293.
- Matteo, F., Charles, L., Harald, P. and Sridhar, R. 1999. Cache-oblivious algorithms. *The 40th Annual Symposium on Foundations of Computer Science* 285–297.
- Min, W. H. and Veeravalli, B. 2005. Aligning biological sequences on distributed bus networks: a divisible load scheduling approach. *IEEE Transactions on Information Technology in Biomedicine* 9 (4): 489–501.
- Moges, M. and Robertazzi, T. G. 2006. Divisible load scheduling and Markov chain models. *Computers and Mathematics with Applications* 52 (10-11): 1529–1542.
- Moges, M., Yu, D. and Robertazzi, T. G. 2009. Grid scheduling divisible loads from two sources. *Computers and Mathematics with Applications* 58 (6): 1081–1092.
- Moges, M. A. and Robertazzi, T. G. 2004. Grid scheduling divisible loads from multiple sources via linear programming. *The 16th International Conference on Parallel and Distributed Computing and Systems (PDCS)* 423–428.
- Montgomery, D. C., Peck, E. A. and Vining, G. G. 2012. *Introduction to linear regression analysis. 5th edition*. John Wiley & Sons.
- Othman, M., Abdullah, M., Ibrahim, H. and Subramaniam, S. 2007. Adaptive divisible load model for scheduling data-intensive grid applications. *Lecture Notes in Computer Science* 4487 (1): 446–453.
- Othman, M., Abdullah, M., Ibrahim, H. and Subramaniam, S. 2008. A^2DLT : divisible load balancing model for scheduling communication-intensive grid applications. *Lecture Notes Computer Science* 5101 (1): 246–253.
- Othman, M., Abdullah, M., Ibrahim, H. and Subramaniam, S. 2009. New optimal load allocation for scheduling divisible data grid applications. *Lecture Notes Computer Science* 5544 (1): 165–174.
- Ottenbacher, Kenneth, J. and Anne, C. 1990. Goal attainment scaling as a method of clinical service evaluation. *American Journal of Occupational Therapy* 44 (6): 519–525.
- Robertazzi, T. G. 1993. Processor equivalence for daisy chain load sharing processors. *IEEE Transactions on Aerospace and Electronic Systems* 29 (4): 1216–1221.

- Robertazzi, T. G. 2003. Ten reasons to use divisible load theory. *Computer* 36 (5): 63–68.
- Robertazzi, T. G. 2005. Divisible load scheduling with multiple sources: closed form solutions .
- Robertazzi, T. G. 2007. *Divisible load modeling for Grids. Networks and Grids Technology and Theory*. 19th edition. New York: Springer.
- Robertazzi, T. G. and Yu, D. 2006. Multi-source grid scheduling for divisible loads. *Proceedings of the 40th Annual Conference on Information Sciences and Systems* 188–191.
- Saaty, T. L. 1988. What is the analytic hierarchy process. *Mathematical Models for Decision Support* 48: 109–121.
- Saaty, T. L. 1990. How to make a decision: the analytic hierarchy process. *European Journal of Operational Research* 48 (1): 9–26.
- Saaty, T. L. 2013. The modern science of multi-criteria decision making and its practical applications: the AHP/ANP approach. *Operation Research* 61 (5): 1101–1118.
- Shi, H., Wang, W. and Kwok, N. 2012. Energy dependent divisible load theory for wireless sensor network workload allocation. *Mathematical Problems in Engineering*. doi:<http://dx.doi.org/10.1155/2012/235289> .
- Shih, C. J. and Chang, C. 1995. Pareto optimization of alternative global criterion method for fuzzy structural design. *Computers & Structures* 54 (3): 455–460.
- Shokripour, A., Othman, M., Ibrahim, H. and Subramaniam, S. 2012. New method for scheduling heterogeneous multi-installment systems. *Future Generation Computer Systems* 28 (8): 1205–1216.
- Sohn, J. and Robertazzi, T. G. 1993. Optimal load sharing for a divisible job on a bus network. *Proceedings of the Conference on Information Sciences and Systems* 835–840.
- Sohn, J. and Robertazzi, T. G. 1998. Optimal time-varying load sharing for divisible loads. *IEEE Transactions on Aerospace and Electronic Systems* 34 (3): 907–923.
- Suresh, S., Huang, H. and Kim, H. J. 2015. Scheduling in compute cloud with multiple data banks using divisible load paradigm. *IEEE Transactions on Aerospace and Electronic Systems* 51 (2): 1288–1297.
- Suresh, S., Kim, H., Run, C. and Robertazzi, T. G. 2012. Scheduling nonlinear divisible loads in a single level tree network. *Journal of Supercomputing* 61 (3): 1068–1088.
- Suresh, S., Mani, V. and Omkar, S. 2003. The effect of start-up delays in scheduling divisible loads on bus networks: an alternate approach. *Computers and Mathematics with Applications* 46 (10): 1545–1557.

- Tudela, A., Akiki, N. and Cisternas, R. 2006. Comparing the output of cost benefit and multi-criteria analysis: An application to urban transport investments. *Transportation Research* 40 (5): 414–423.
- V., K. V., Moscicki, J. T. and Krzhizhanovskaya, V. 2009. Dynamic workload balancing of parallel applications with user-level scheduling on the grid. *Future Generation Computer Systems* 25 (1): 28–34.
- Veeravalli, B. and Barlas, G. 2000. Access time minimization for distributed multimedia applications. *Multimedia Tools and Applications* 12 (2-3): 235–256.
- Veeravalli, B. and Gerassimos, B. 2003. Scheduling divisible loads with processor release times and finite size buffer capacity constraints in bus networks. *Cluster Computing* 6 (1): 63–74.
- Veeravalli, B., Ghose, D. and Mani, V. 1995. Multi-installment load distribution in tree networks with delays. *IEEE Transactions on Aerospace and Electronic Systems* 31 (2): 555–567.
- Veeravalli, B., Ghose, D., Mani, V. and Robertazzi, T. G. 1996. *Scheduling divisible loads in parallel and distributed systems*. 8th edition. IEEE.
- Veeravalli, B., Ghose, D. and Robertazzi, T. G. 2003. Divisible load theory: A new paradigm for load scheduling in distributed systems. *Cluster Computing* 6 (1): 7–17.
- Veeravalli, B., Li, X. and Ko, C. C. 2000a. Design and analysis of load distribution strategies with start-up costs in scheduling divisible loads on distributed networks. *Mathematical and Computer Modeling* 32 (7): 901–932.
- Veeravalli, B., Li, X. and Ko, C. C. 2000b. Efficient partitioning and scheduling of computer vision and image processing data on bus networks using divisible load analysis. *Image and Vision Computing* 18 (11): 919–938.
- Veeravalli, B. and Ranganath, S. 2002. Theoretical and experimental study on large size image processing applications using divisible load paradigm on distributed bus networks. *Image and Vision Computing* 20 (13): 917–935.
- W., G. 1998. A Greedy Algorithm for Processing a Divisible Load on a Hypercube. *International Conference on Parallel Computing in Electrical Engineering PARELEC98* 1: 185–188.
- Wolniewicz, P. 2002. Multi-installment divisible job processing with communication start up cost. *Foundations of Computing and Decision Sciences* 27 (1): 43–57.
- Wolniewicz, P. 2003. *Divisible job scheduling in systems with limited memory*. Ph.D. dissertation. Poznan University of Technology, Poland.
- Wolniewicz, P. and Drozdowski, M. 2002. Processing time and memory requirements for multi-installment divisible job processing. *Lecture Notes in Computer Science* 2328: 125–133.

- Wong, H. M., Yu, D., Veeravalli, B. and Robertazzi, T. G. 2003. Data intensive grid scheduling: multiple sources with capacity constraints. *Proceedings of the 15th International Conference on Parallel and Distributed Computing and Systems* 1: 7–11.
- Woo, L. J. and Hie, K. S. 2000. Using analytic network process and goal programming for interdependent information system project selection. *Computers & Operations Research* 27 (4): 367–382.
- Yang, Y. and Casanova, H. 2003. Rumr: Robust scheduling for divisible workloads. *Proceedings of the 12th IEEE International Symposium on High Performance Distributed Computing* 114–123.
- Yao, J. and Veeravalli, B. 2004. Design and performance analysis of divisible load scheduling strategies on arbitrary graphs. *Cluster Computing* 7 (2): 191–207.
- Yuan-Chieh, C. and Robertazzi, T. G. 1988. Distributed computation with communication delay (distributed intelligent sensor networks). *IEEE Transactions on Aerospace and Electronic Systems* 24 (6): 700–712.