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The chromaticity of s-bridge graphs and related graphs

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Abstract

The graph consisting of s paths joining two vertices is called an s-bridge graph. In this paper, we discuss the chromaticity of some families of s-bridge graphs, especially 4-bridge graphs, and some graphs related to s-bridge graphs.

1. The chromaticity of 4-bridge graphs

The graphs considered here are finite, undirected, simple and loopless. For a graph G, let V(G) denote the vertex set of G and E(G) the edge set of G. Let $P(G; \lambda)$, or simply P(G) if there is no likelihood of confusion, denote the chromatic polynomial of G. In this paper, $y = \lambda - 1$ and x = -y. Two graphs G and H are said to be chromatically equivalent if P(G) = P(H). A graph G is said to be chromatically unique if P(G) = P(H) implies H is isomorphic to G. Let $\mathcal{G} = \{G_0, \ldots, G_p\}$ and $\mathcal{K} = \{K_{i_1}, \ldots, K_{i_p}\}$. Then all graphs obtained from G_0, \ldots, G_p by overlapping in K_{i_1}, \ldots, K_{i_p} in different positions and different orders form a class of graphs. We denote it by $\{\mathcal{G}, \mathcal{K}\}$. If the chromatic equivalence of H to $\{\mathcal{G}, \mathcal{K}\}$ implies $H \in \{\mathcal{G}, \mathcal{K}\}$, then $\{\mathcal{G}, \mathcal{K}\}$ is said to be a complete class of chromatically equivalent graphs. For details and other symbols and definitions, readers can see [6].

A 2-connected graph G is called a generalized polygon tree if it can be decomposed into cycle class $\mathcal{G} = \{C_{i_1}, \ldots, C_{i_r}\}$ and there exist an overlapping process: $H_1 = C_{i_1}, H_j$ is obtained from H_{j-1} and C_{i_j} by overlapping in path P_{i_j} where in each step of overlapping, the vertices on P_{i_j} , except end vertices, are with degree 2.

In [6], it was proved that a 2-connected graph is a generalized polygon tree if and only if it has no subgraphs homeomorphic to K_4 . Obviously a generalized polygon

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tree is a planar graph. For a 2-connected planar graph G, we define r(G) as the number of interior regions of G: r(G) = |E(G)| - |V(G)| + 1. For a generalized polygon tree, we define intercourse number of G, $\sigma(G)$, as the number of nonadjacent vertex pairs, where there are at least three internally disjoint paths joining them. It was also proved in [6] that if G is a generalized polygon tree and P(H) = P(G), then H is also a generalized polygon tree and r(H) = r(G), $\sigma(H) = \sigma(G)$.

A graph consisting of s paths joining two vertices is called an s-bridge graph, which is denoted by $F(k_1, ..., k_s)$, where $k_1, ..., k_s$ are the lengths of s paths. Clearly an s-bridge graph is a generalized polygon tree.

It was proved by Chao and Whitehead [1] that the cycle C_n is chromatically unique. Later, Loerinc [3] proved that the generalized θ -graph θ_{abc} is chromatically unique. That is to say, 2-bridge graphs and 3-bridge graphs are all chromatically unique. In this paper, we consider the chromaticity of s-bridge graphs and some graphs related to s-bridge graphs. First we give the sufficient and necessary conditions for a 4-bridge graph to be chromatically unique.

Now let G be a 4-bridge graph F(a,b,c,d), the lengths of the 4 paths joining two vertices u and v be a, b, c and d, where $a \ge b \ge c \ge d$ and |V(G)| = n. Then |E(G)| = a + b + c + d = n + 2.

Theorem 1. A 4-bridge graph G is not chromatically unique if and only if G satisfies one of the following conditions:

- (1) d = 1,
- (2) d=2 and a=b+1=c+2.

Proof. If d=1, G is a polygon tree obtained from C_{a+1} , C_{b+1} and C_{c+1} by overlapping on an edge. Let $\mathcal{G} = \{C_{a+1}, C_{b+1}, C_{c+1}\}$, $\mathcal{K} = \{K_2, K_2\}$. As shown in [6], $\{\mathcal{G}, \mathcal{K}\}$ is a complete class of chromatically equivalent graphs and it is easy to see that $|\{\mathcal{G}, \mathcal{K}\}| > 1$ (see Fig. 1.) Hence G is not chromatically unique and all graphs which are chromatically equivalent to G belong to $\{\mathcal{G}, \mathcal{K}\}$.

In the following, we always assume that d>1. Then G is a generalized polygon tree, the interior region number r(G)=3 and the intercourse number $\sigma(G)=1$. First, using the formula P(G)=P(G+uv)+P(G.uv), we compute the chromatic





Fig. 1.

polynomial of G.

$$\begin{split} P(G) &= \frac{1}{x^3(x-1)^3} \, P(C_{a+1}) \, P(C_{b+1}) P(C_{c+1}) P(C_{d+1}) \\ &- \frac{1}{(x-1)^3} \, P(C_a) P(C_b) P(C_c) P(C_d) \\ &= \frac{(-1)^{n-1} x}{(x-1)^2} \, \left[(1+x+x^2) - (x+1)(x^a+x^b+x^c+x^d) \right. \\ &+ (x^{a+b}+x^{a+c}+x^{a+d}+x^{b+c}+x^{b+d}+x^{c+d}) - x^{n+1} \right] \\ &= \frac{(-1)^{n-1} x}{(x-1)^2} \, Q(G). \end{split}$$

Suppose that there is a graph H such that P(G) = P(H). By Lemma 4 in [6], we know that H is also a generalized polygon tree and the interior region number r(H) = r(G) = 3, the intercourse number $\sigma(H) = \sigma(G) = 1$, i.e. H is either a 4-bridge graph or a graph obtained from a generalized θ -graph and a cycle by overlapping on an edge.

Assume that H is a 4-bridge graph with a', b', c' and d' to be the lengths of its paths, where $a' \ge b' \ge c' \ge d'$. Comparing the coefficients of the terms with the lowest degrees in Q(G) and Q(H), we conclude that d' must be equal to d. By Lemma 4 in [6], the girth g(H) = g(G), hence we have c' = c. It is easy to obtain that b' = b and a' = a. Therefore, H is isomorphic to G.

Now suppose that H is obtained from a generalized θ -graph $\theta_{a'b'c'}$ and a cycle $C_{d'}$ by overlapping on an edge, where $a' \ge b \ge c' \ge 2$ and $d' \ge 3$, a' + b' + c' + d' - 3 = |V(H)| = n, i.e. a' + b' + c' + d' = a + b + c + d + 1. Since the girth of G is more than 3, the girth of G is also more than 3, therefore $d' \ge 4$. We compute the chromatic polynomial of G.

$$P(H) = \frac{1}{x(x-1)} P(\theta_{a'b'c'}) P(C_{a'})$$

$$= \frac{(-1)^{n-1}x}{(x-1)^2} \left[(1+x) - (x^{a'-1} + x^{a'} + x^{c'} + x^{b'} + x^{a'}) + (x^{c'+a'-1} + x^{b'+a'-1} + x^{a'+a'-1}) + x^{a'+b'+c'-1} - x^{n+1} \right]$$

$$= \frac{(-1)^{n-1}x}{(x-1)^2} Q(H).$$

Now we solve the equation Q(G) = Q(H). After canceling x^{n+1} , x and constant terms, we get

$$a+b+c+d=a'+b'+c'+d'-1$$
, $a \ge b \ge c \ge d \ge 2$, $a' \ge b' \ge c' \ge 2$, $a' \ge 4$.

$$Q(G): \quad x^{2} - x^{d} - x^{d+1} - x^{c} - x^{c+1} - x^{b} - x^{b+1} - x^{a} - x^{a+1}$$

$$+ x^{c+d} + x^{b+d} + x^{a+d} + x^{b+c} + x^{a+c} + x^{a+b}.$$

$$Q(H): \quad -x^{d'-1} - x^{d'} - x^{c'} - x^{b'} - x^{a'} + x^{c'+d'-1}$$

$$+ x^{b'+d'-1} + x^{a'+d'-1} + x^{a'+b'+c'-1}$$

Since $d = \min\{d, d+1, c, c+1, b, b+1, a, a+1\}$ and $\min\{a'+d'-1, b'+d'-1, c'+d'-1, a'+b'+c'-1\} > 2$, x^2 of Q(G) can only be canceled with x^d of Q(G), i.e. d=2.

$$a+b+c+3 = a'+b'+c'+d', \ a \ge b \ge c \ge 2, \ a' \ge b' \ge c' \ge 2, \ d' \ge 4.$$

$$Q(G): -x^3 - x^c - x^{c+1} - x^b - x^{b+1} - x^a - x^{a+1} + x^{c+2} + x^{b+2} + x^{a+2} + x^{b+c} + x^{a+c} + x^{a+b}$$

$$Q(H): -x^{d'-1} - x^{d'} - x^{c'} - x^{b'} - x^{a'} + x^{c'+d'-1}$$

$$+ x^{b'+d'-1} + x^{a'+d'-1} + x^{a'+b'+c'-1}$$

Since $\min\{c+2, b+2, a+2, b+c, a+c, a+b\} > 3$, $-x^3$ of Q(G) can only be canceled with one term of Q(H), i.e. $-x^{d'-1}$, $-x^{d'}$, $-x^{c'}$, $-x^{b'}$ or $-x^{a'}$. Because $d' \ge 4$, $-x^3$ cannot be canceled with $-x^{d'}$. If $-x^3$ is canceled with $-x^{d'-1}$, then d'=4; otherwise one of c', b' and a' must be equal to 3. Without loss of generality we may assume d'=4 or c'=3. Thus we consider the following two cases.

Case 1. d' = 4.

In this case, we have

$$a+b+c=a'+b'+c'+1, \quad a\geqslant b\geqslant c\geqslant 2, \quad a'\geqslant b'\geqslant c'\geqslant 2.$$

$$Q(G): \quad -x^{c}-x^{c+1}-x^{b}-x^{b+1}-x^{a}-x^{a+1}$$

$$+x^{c+2}+x^{b+2}+x^{a+2}+x^{b+c}+x^{a+c}+x^{a+b}$$

$$Q(H): \quad -x^{4}-x^{c'}-x^{b'}-x^{a'}+x^{c'+3}+x^{b'+3}+x^{a'+3}+x^{a'+b'+c'-1}$$

By observing the terms with the highest power, we may conclude that a+b=a'+b'+c'-1 and c=2. Thus c'=2. We now have

$$a+b=a'+b'+1$$
, $a \ge b \ge 2$, $a' \ge b' \ge 2$.
 $Q(G)$: $-x^3-x^b-x^{b+1}-x^a-x^{a+1}+x^4+2x^{b+2}+2x^{a+2}$.
 $Q(H)$: $-x^4-x^{b'}-x^{a'}+x^5+x^{b'+3}+x^{a'+3}$.

Then we have a = a' + 1 = b' + 1 = b + 1.

$$Q(G)$$
: $-x^3-2x^{b+1}+x^4+x^{b+2}$.

$$Q(H)$$
: $-x^4-x^b+x^5$.

The solution is b=3. Therefore F(4,3,2,2) is chromatic equivalent with the graph obtained from F(3,3,2) and C_4 by overlapping on an edge.

Case 2.
$$c' = 3$$
.

In this case we have

$$a+b+c=a'+b'+d', \quad a\geqslant b\geqslant c\geqslant 2, \quad a'\geqslant b'\geqslant 2, \quad d'\geqslant 5.$$

$$Q(G): \quad -x^{c}-x^{c+1}-x^{b}-x^{b+1}-x^{a}-x^{a+1}$$

$$+x^{c+2}+x^{b+2}+x^{a+2}+x^{b+c}+x^{a+c}+x^{a+b},$$

$$Q(H): \quad -x^{d'-1}-x^{d'}-x^{b'}-x^{a'}+x^{d'+2}+x^{b'+d'-1}+x^{a'+d'-1}+x^{a'+b'+2}.$$

By comparing the least power terms, we can see that c=d'-1 or c=b'. If c=d'-1, then the girth of G is c+2>d' which is a contradiction. So c=b'. We now have

$$a+b=a'+d', a \geqslant b \geqslant c, a' \geqslant c \geqslant 2, d' \geqslant 5.$$

$$Q(G): -x^{c+1} - x^b - x^{b+1} - x^a - x^{a+1} + x^{c+2} + x^{b+2} + x^{a+2} + x^{b+c} + x^{a+c} + x^{a+b}.$$

$$Q(H): -x^{d'-1} - x^{d'} - x^{a'} + x^{d'+2} + x^{c+d'-1} + x^{a'+d'-1} + x^{a'+c+2}.$$

We now compare the terms of highest power. Since $c+d'-1 \le a'+d'-1 = a+b-1$, we have a+b=a'+c+2=a'+d'. Thus d'=c+2. So we get

$$a+b=a'+c+2$$
, $a \ge b \ge c$, $a' \ge c \ge 3$.
 $Q(G)$: $-x^b-x^{b+1}-x^a-x^{a+1}+x^{c+2}+x^{b+2}+x^{a+2}+x^{b+c}+x^{a+c}$.
 $Q(H)$: $-x^{c+2}-x^{a'}+x^{c+4}+x^{2c+1}+x^{a'+c+1}$.

In order that x^{c+2} can be canceled, at least two of a, b, a+1 and b+1 must be equal to c+2. If a=b=c+2, then a'=c+2. There is no solution. If a=b=c+1, then a'=c. Again, no solution. If a=b+1=a+2, then a'=a+1 and we have the solution. Thus F(c+2,c+1,c,2) is chromatic equivalent with the graph obtained from F(c+1,c,3) and C_{c+2} by overlapping on an edge.

Now we have solved the equation Q(G) = Q(H) and completed the proof. From the above argument we can see that when a = b + 1 = c + 2 and d = 2, the graphs obtained from $\theta_{3,c,c+1}$ and C_{c+2} by overlapping on an edge are the only graphs chromatically equivalent to F(c+2,c+1,c,2). \square

2. Some complete classes of chromatically equivalent graphs

Now we consider the chromaticity of some s-bridge graphs and graphs related to s-bridge graphs. First we consider the class $\{\{\theta_{abc}, C_d\}, \{K_2\}\}\}$, i.e. the set of those graphs that are obtained from a θ -graph and a cycle by overlapping on an edge.

Theorem 2. $\{\{\theta_{abc}, C_d\}, \{K_2\}\}\$, where $a \ge b \ge c \ge 2$ and $d \ge 3$, is not a complete class of chromatically equivalent graphs if and only if c = 3 and d = a + 1 = b + 2.

Proof. Let $G \in \{\{\theta_{abc}, C_d\}, \{K_2\}\}\}$. Then G is a generalized polygon tree, r(G) = 3 and $\sigma(G) = 1$. By Lemma 4 in [7], if there is a graph H such that P(G) = P(H), then H is either a 4-bridge graph or a graph obtained from a generalized θ -graph and a cycle by overlapping on an edge. If H is a 4-bridge graph, and the lengths of whose 4 paths are a', b', c' and d', then by Theorem 1, P(G) = P(H) if and only if c = 3, d = a + 1 = b + 2, d' = 2, a' = b' + 1 = c' + 2 and b = c'. So we only need to prove that if $H \in \{\{\theta_{a'b'c'}, C_{d'}\}, \{K_2\}\}$, where $a' \geqslant b' \geqslant c'$, and P(G) = P(H), then a' = a, b' = b, c' = c and d' = d. Let a + b + c + d = a' + b' + c' + d' = n. If d' = d, then by the chromatic uniqueness of the generalized θ -graph, it follows a' = a, b' = b and c' = c. So in the following we prove that if $d \neq d'$, then there is no solution for P(G) = P(H).

As done in Section 1, let $P(G) = [(-1)^{n-1}x/(x-1)^2]Q(G)$ and $P(H) = [(-1)^{n-1}x/(x-1)^2]Q(H)$. We try to solve the equation Q(G) = Q(H). Canceling some terms, we get

$$a+b+c+d=a'+b'+c'+d', \quad a \ge b \ge c \ge 2, \quad a' \ge b' \ge c' \ge 2, \quad d \ne d'.$$

$$Q(G): \quad -x^{d-1}-x^d-x^c-x^b-x^a+x^{c+d-1}+x^{b+d-1}+x^{a+d-1}+x^{a+b+c-1}.$$

$$Q(H): \quad -x^{d'-1}-x^{d'}-x^{c'}-x^{b'}-x^{a'}+x^{c'+d'-1}+x^{b'+d'-1}+x^{a'+d'-1}+x^{a'+b'+c'-1}$$

Without loss of generality, we can assume d > d'. Then the girth of G must be b + c satisfying d > b + c and the girth of H is either d' or b' + c' if b' + c' < d'. In both cases, comparing the terms with lowest degrees of Q(G) and Q(H) yields c' = c and b' = b.

$$Q(G): -x^{d-1} - x^d - x^a + x^{c+d-1} + x^{a+b+c-1}.$$

$$Q(H): -x^{d'-1} - x^{d'} - x^{a'} + x^{c'+d'-1} + x^{a'+b'+c'-1}.$$

The solutions are (1) a' = a and d' = d; (2) a' = d - 1, a = d' - 1 and c = c' = 1, where the first one contradicts the assumption that d > d' and the second one contradicts the condition that $c \ge 2$. Thus we have finished the proof. \Box

Since $|\{\{\theta_{2,2,2},C_d\},\{K_2\}\}|=1$, the graph obtained from $\theta_{2,2,2}$ and a cycle by overlapping on an edge is chromatically unique.

Theorem 3. Let $\mathcal{G} = \{G_0, ..., G_k\}$ and $\mathcal{K} = k\{K_2\}$, where $G_0, ..., G_k$ are all nonseparable generalized polygon trees. Then, if $\{\mathcal{G}, \mathcal{K}\}$ is a complete class of chromatically equivalent graphs, so is $\{\mathcal{G} \cup s\{K_3\}, \mathcal{K} \cup s\{K_2\}\}$ for any positive integer s.

Proof. We first prove that if a graph H is a nonseparable generalized polygon tree with r(H) > 1, then g(H) > 3. Now H can be obtained from some cycles by overlapping on paths. If one cycle has length 3, it must overlap on edges with other subgraphs, i.e. H is separable. A contradiction.

Suppose that G is chromatically equivalent to $\{\mathscr{G} \cup s\{K_3\}, \mathscr{K} \cup s\{K_2\}\}$. Then G is a generalized polygon tree. Let $G \in \{\mathscr{G}', \mathscr{K}'\}$. Then all graphs in \mathscr{G}' must be generalized polygon trees. Let $\mathscr{G}' = \{H_0, \ldots, H_t\}$ and $\mathscr{K}' = t\{K_2\}$, where H_0, \ldots, H_t are all nonseparable. If $r(H_i) > 1$, then $g(H_i) > 3$. Therefore \mathscr{G}' must contain at least s triangles. So $\{\mathscr{G}' \setminus s\{K_3\}, \mathscr{K}' \setminus s\{K_2\}\}$ is chromatically equivalent to $\{\mathscr{G}, \mathscr{K}\}$. Since $\{\mathscr{G}, \mathscr{K}\}$ is a complete class of chromatically equivalent graphs, $\{\mathscr{G}' \setminus s\{K_3\}, \mathscr{K}' \setminus s\{K_2\}\} = \{\mathscr{G}, \mathscr{K}\}$, i.e. $\{\mathscr{G}', \mathscr{K}'\} = \{\mathscr{G} \cup s\{K_3\}, \mathscr{K} \cup s\{K_2\}\}$. Therefore $\{\mathscr{G} \cup s\{K_3\}, \mathscr{K} \cup s\{K_2\}\}$ is a complete class of chromatically equivalent graphs. \square

Theorem 4. Let G be an s-bridge graph, where the lengths of the s paths are all m. Then G is chromatically unique.

Proof. The vertex number of G is (n-1)s+2, g(G)=2n, the number of cycles in G with length 2n is $\binom{s}{2}$, and G is a generalized polygon tree with r(G)=s-1 and $\sigma(G)=1$.

Now if there is a graph H such that P(H) = P(G), then |V(H)| = (n-1)s + 2, g(H) = 2n and H is also a generalized polygon tree with $\sigma(H) = 1$. So H is obtained from a t-bridge graph H' and s-t cycles by overlapping on edges where $t \le s$. Since g(H) = 2n, if one path of t-bridge graph H' has length n' < n, the others must have lengths at least 2n-n'. So $|V(H')| \ge (n-1)t+2$ and the number of cycles in H' with length 2n is at most $\binom{t}{2}$. It is easy to see that the number of cycles in H with length 2n is at most $\binom{t}{2} + (n-1)(s-t)/(2n-2) \le \binom{s}{2}$ and the equality holds if and only if t = s and H is isomorphic to G. Hence G is chromatically unique. \square

As a special case, $K_{2,s}$ is chromatically unique, which was proved in [4] using other methods.

Theorem 5. Let G be an s-bridge graph, the lengths of whose s paths satisfy $j_1 \ge \cdots \ge j_s \ge s-1$. Then G is chromatically unique.

Proof. Suppose that there is a graph H such that P(H) = P(G). Then H is either an s-bridge graph or a graph obtained from a t-bridge graph and s-t cycles by overlapping on edges, where t < s.

If H is an s-bridge graph, we solve equation P(H) = P(G) and it is easy to conclude that H is isomorphic to G.

Suppose that H is obtained from a t-bridge graph and s-t cycles by overlapping on edges, where the lengths of t paths in t-bridge graph are $k_1 \ge \cdots \ge k_t$, and the lengths of s-t cycles are $l_1 \ge \cdots \ge l_{s-t}$. First we give the chromatic polynomials

of G and H.

$$P(G) = \frac{y}{(y+1)^{s-1}} \prod (y^{j_i} + (-1)^{j_i+1}) + \frac{y^s}{(y+1)^{s-1}} \prod (y^{j_i-1} + (-1)^{j_i})$$

$$= \frac{y}{(y+1)^{s-1}} \left(\prod (y^{j_i} + (-1)^{j_i+1}) + y^{s-1} \prod (y^{j_i-1} + (-1)^{j_i}) \right)$$

$$= \frac{y}{(y+1)^{s-1}} Q(G),$$

$$P(H) = \frac{y}{(y+1)^{s-1}} \left(\prod (y^{k_i} + (-1)^{k_i+1}) + y^{t-1} \prod (y^{k_i-1} + (-1)^{k_i}) \right)$$

$$= \frac{y}{(y+1)^{s-1}} Q(H).$$

We solve the equation Q(G) = Q(H). Since |V(G)| = |V(H)|, we have $\sum j_i - s + 2 = \sum k_i + \sum l_i - 2s + t + 2$. It is easy to see that the term with the lowest power in

$$Q(G)-(-1)^{\sum j_i+s}$$

must be either

$$(-1)^{\sum j_i} y^{s-1}$$
 or $(-1)^{\sum_{i\neq s} j_i - s + 1} y^{j_s}$,

which cannot be canceled with each other. The corresponding term in

$$O(H)-(-1)^{\sum k_i+\sum l_i+t}$$

must be one of the three terms:

$$(-1)^{\sum_{i\neq i}k_i+\sum l_i-t+1}y^{k_i}, (-1)^{\sum k_i+\sum l_i}y^{t-1}$$
 and $(-1)^{\sum_{i\neq s-t}l_i+\sum k_i-t}y^{l_{s-t}-1},$

which cannot be canceled either, so

$$\min\{t-1, k_t, l_{s-t}-1\} = \min\{s-1, j_s\}.$$

Since $t-1 < s-1 \le j_s$, the equality cannot hold, i.e. there is no solution for P(G) = P(H). \square

Theorem 6. Let M be an s-bridge graph, the lengths of whose s paths are all n. Then $\{\{M, C_p\}, \{K_2\}\}\$ is a complete class of chromatically equivalent graphs.

Proof. First we give the chromatic polynomial of $G \in \{\{M, C_p\}, \{K_2\}\}$ as follows:

$$P(G) = \frac{y}{(y+1)^s} ((y^n + (-1)^{n+1})^s + y^{s-1} (y^{n-1} + (-1)^n)^s) (y^{p-1} + (-1)^p)$$
$$= \frac{y}{(y+1)^s} Q(G).$$

The term of $Q(G)-(-1)^{ns+s+p}$ with the lowest power must be one of the three terms: $(-1)^{(s-1)(n+1)+p}sy^n$, $(-1)^{sn+p}y^{s-1}$ and $(-1)^{sn+s}y^{p-1}$ which cannot be canceled with each other.

Suppose that there is a graph H such that P(H) = P(G). Then H is either an (s+1)-bridge graph or a graph obtained from a t-bridge graph $(t \le s)$ and cycles by overlapping on edges.

Case 1. H is an (s+1)-bridge graph. If $g(G) = 2n \le p$, then because G has $\binom{s}{2}$ or $\binom{s}{2}+1$ cycles with length 2n, it is easy to conclude that there is no solution for P(G) = P(H). So p < 2n. Let the lengths of s+1 paths of H be $j_1 \ge \cdots \ge j_{s+1}$. Then $j_s + j_{s+1} = p$ and

$$P(H) = \frac{y}{(y+1)^s} \left(\prod (y^{j_i} + (-1)^{j_i+1}) + y^s \prod (y^{j_i-1} + (-1)^{j_i}) \right)$$
$$= \frac{y}{(y+1)^s} Q(H).$$

The term in

$$Q(H) - (-1)^{\sum j_i + s + 1}$$

with the lowest power is either

$$(-1)^{\sum_{i\neq s+1}j_i+s}v^{j_{s+1}}$$
 or $(-1)^{\sum j_i}v^s$,

which cannot be canceled with each other. So

$$\min\{i_{s+1}, s\} = \min\{s-1, n, p-1\}.$$

We can only get $j_{s+1} = s-1$ and s = n. However, there is no solution for Q(G) = Q(H). Case 2. H is obtained from a t-bridge graph and the cycles $C_{l_1}, \ldots, C_{l_{s-t+1}}$ by overlapping on edges, where t < s, the lengths of the t paths of the t-bridge graph are $j_i \ge \cdots \ge j_t$ and $l_1 \ge \cdots \ge l_{s-t+1}$.

If the equality $l_k = p$ holds for an integer k, we can know that $H \in \{\{M, C_p\}, \{K_2\}\}\}$. If $2n \le p$, since g(G) = 2n and G has $\binom{s}{2}$ or $\binom{s}{2} + 1$ cycles with length 2n, it is easy to conclude that $H \in \{\{M, C_p\}, \{K_2\}\}\}$. Now g(G) = p < 2n, $j_t + j_{t-1} = p < l_{s-t+1}$ and

$$P(H) = \frac{y}{(y+1)^{s}} \left(\prod (y^{j_{i}} + (-1)^{j_{i}+1}) + y^{t-1} \prod (y^{j_{i}-1} + (-1)^{j_{i}}) \right)$$

$$\prod (y^{l_{i}-1} + (-1)^{l_{i}})$$

$$= \frac{y}{(y+1)^{s}} Q(H).$$

Then

$$\min\{p-1,n,s-1\} = \min\{l_{s-t+1}-1,t-1,j_t\}.$$

Since t < s, $j_t < n$ and $j_t < p-1$, there is no solution for Q(G) = Q(H).

Since $|\{\{K_{2,s}, C_p\}, \{K_2\}\}| = 1$ the graph obtained from $K_{2,s}$ and C_p by overlapping on an edge is chromatically unique.

At last, we give the following theorem without proof, which is similar to Theorems 5 and 6.

Theorem 7. Let M be an s-bridge graph with s paths whose lengths are $j_1 \ge \cdots \ge j_s > s$. Then $\{\{M, C_p\}, \{K_2\}\}$ is a complete class of chromatically equivalent graphs.

Now if M is an s-bridge graph, then $|\{\{M, C_p\}, \{K_2\}\}\}| = ?$ Let $\lceil a \rceil$ be the minimum integer equal to or more than a. The following result is obvious: Suppose that the s-bridge graph M has t paths with different lengths $j_1 > \cdots > j_t$. Then

$$|\{\{M, C_p\}, \{K_2\}\}| = \sum_{i=1}^t \lceil j_i/2 \rceil.$$

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